

## Lecture 3: August 23

Lecturer: Neil Calkin

Scribe: Jeannie Friedel and Dominique Morgan

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### 3.1 Combinatorial Classes:

We can extend the cross product of two combinatorial classes to finitely many  $A \times B \times \dots \times C \times D$ .

If  $A, B, \dots, C, D$  respectively have generating functions  $f_a, f_b, \dots, f_c, f_d$ , then  $A \times B \times \dots \times C \times D$  has generating function  $f_a(x) f_b(x) \dots f_c(x) f_d(x)$ , where the weight is  $w(a, b, \dots, c, d) = w_a(a) + w_b(b) + \dots + w_c(c) + w_d(d)$ .

#### 3.1.1 Example

Let  $A = \{0, 1\}$  with  $w(0) = 0$  and  $w(1) = 1$  so  $A$  has generating function  $1 + x$ .

Let  $A^n$  denote  $\underbrace{A \times A \times \dots \times A}_{n\text{-times}}$  which has generating function  $(1+x)^n$ .

A typical element of  $A^n$  is an  $n$ -tuple of 0's and 1's (a binary sequence). Equivalently, a sequence of 0's and 1's of length  $n$ . A sequence has weight  $k$  if it has exactly  $k$  1's.

Bijection  $\phi_n : \{\text{Set of binary strings of length } n\} \longrightarrow \{\text{Set of subsets of } \{1, 2, \dots, n\}\}$ .

$|\phi_n(\sigma)| = w(\sigma)$  and  $\phi_n = \{j \mid \sigma(j) = 1\} = \text{set of positions in which } \sigma \text{ has a } 1$ .

e.g. 0110110111  $\rightarrow \{2, 3, 5, 6, 8, 9, 10\}$

#### Corollary 3.1

Number of subsets of  $\{1, 2, \dots, n\}$  of cardinality  $k$   
 = Number of binary strings of length  $n$  with  $k$  1's  
 = Number of binary strings of length  $n$  with weight  $k$   
 =  $[x^k] (1+x)^n$

Define  $n \geq k \geq 0$ ,  $n, k$  integer,  $\binom{n}{k} = [x^k] (1+x)^n$ .

**Corollary 3.2**  $\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}$  *Pascal's identity*

**Proof:**

$$\begin{aligned} [x^k](1+x)^n &= [x^k](1+x)^{n-1} + x(1+x)^{n-1} = \left([x^k](1+x)^{n-1}\right) + \left([x^k]x(1+x)^{n-1}\right) \\ &= \binom{n-1}{k} + [x^{k-1}](1+x)^{n-1} = \binom{n-1}{k} + \binom{n-1}{k-1} \end{aligned}$$

**Proof:** [Combinatorial]

Let  $\binom{n}{k}$  be the number of  $k$ -subsets of  $\{1, 2, \dots, n\} = |S_{n,k}|$  where  $S_{n,k}$  is the set of  $k$ -subsets of  $\{1, 2, \dots, n\}$ . Let  $T_{n,k}$  be the set of  $k$ -subsets of  $\{1, 2, \dots, n\}$  which contain the element  $n$ .

$S_{n,k} = S_{n-1,k} \dot{\cup} T_{n,k}$  where  $\dot{\cup}$  is the disjoint union, that is  $S_{n-1,k} \cap T_{n,k} = \emptyset$ .

Now,  $|T_{n,k}| = |S_{n-1,k-1}|$ , indeed there is a natural bijection (delete the element  $n$ ). Hence  $|S_{n,k}| = |S_{n-1,k}| + |S_{n-1,k-1}|$ , and so

$$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}$$

as required. ■

### 3.1.2 Sets of Binary Strings via regular languages:

Take a nice language  $S$ ; denote by  $S^*$  the set of all words formed by concatenating a finite number of elements of  $S$ . “nice” means that each word is obtained uniquely.

Example:  $S = \{0, 1\}$ .  $S^* = \{\epsilon, 0, 1, 00, 01, 10, 11, 000, 001, 010, 011, 100, 110, 111, \dots\}$ , where  $\epsilon$  is the empty string.

Bad Example: The set  $\{0, 1, 01\}$  is not nice since the string  $01$  is obtained as  $01$  and also  $(0)(1)$ , so not obtained uniquely.

Is the set  $\{0, 01\}$  nice? yes. We can represent it the following ways:

$$0^{a_1} \dot{\cup} 0^{a_1} 10^{a_2} \dot{\cup} 0^{a_1} 10^{a_2} 10^{a_3} \dot{\cup} \dots$$

$$0^{b_1} \dot{\cup} 0^{b_1} (01)0^{b_2} \dot{\cup} 0^{b_1} (01)0^{b_2} (01)0^{b_3} \dot{\cup} \dots$$

Back to considering binary sequences.  $S = \{0, 1\}$ , and

$$\begin{aligned} S^* &= \text{all binary sequences} \\ &= \{0\}^* \{1\{1\}^* 0\{0\}^*\}^* \{1\}^* \\ &\quad \text{which we'll write} \\ &= 0^*(11^*00^*)^*1^* \end{aligned}$$

What about the set of binary strings without consecutive 1's?

$$\{\epsilon, 0, 1, 00, 01, 10, 000, 001, 010, 100, 101, \dots\} = 0^*(100^*)^*(\epsilon \cup 1)$$

Let's see if we can compute the generating function for such strings with  $x$  marking length.

$$1 + 2x + 3x^2 + 5x^3 + \dots$$

$0^*$  has generating function  $1 + x + x^2 + x^3 + \dots = \frac{1}{1-x}$ .  $100^*$  has generating function  $\frac{x^2}{1-x}$ . So,  $(100^*)^*$  has generating function

$$1 + \left(\frac{x^2}{1-x}\right) + \left(\frac{x^2}{1-x}\right)^2 + \left(\frac{x^2}{1-x}\right)^3 + \dots = \frac{1}{1 - \frac{x^2}{1-x}}$$

$\epsilon \cup 1$  has generating function  $1 + x$ . So,  $0^*(100^*)^*(\epsilon \cup 1)$  has generating function

$$\frac{1}{1-x} \cdot \frac{1}{1 - \frac{x^2}{1-x}} \cdot (1+x) = \frac{1+x}{1-x-x^2}.$$

How can we get the Fibonacci numbers from this???