

## Lecture 8: September 3

Lecturer: Neil Calkin

Scribe: Erin Doolittle and Jeannie Friedel

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## 8.1 Power Series Inverse

Example:

$$(1+x)(1-x+x^2-x^3+x^4+\dots)$$

converges in the ring of formal power series. So we write

$$(1+x)^{-1} = 1 - x + x^2 - x^3 + x^4 + \dots$$

Since we are in a commutative ring (for now) so  $f(x)g(x) = g(x)f(x)$ . If  $f(x)$  has an inverse,  $g(x)$  say, that is  $f(x)g(x) = 1$  then it is unique.

(Non-commutative version: if  $f(x)$  has a right inverse  $g(x)$  and a left inverse  $h(x)$  then  $g(x) = h(x)$ .)

**Proof:**

$$\begin{aligned} h(x)f(x) &= 1 & f(x)g(x) &= 1 \\ g(x) &= (h(x)f(x))g(x) = h(x)(f(x)g(x)) = h(x) \end{aligned}$$

■

## 8.2 Which power series have an inverse?

Suppose  $f(x)$  has coefficient in a commutative ring,  $R$  (e.g.  $\mathbb{Z}$ ). When does there exist  $g(x)$  power series over the same ring so that  $f(x)g(x) = 1$ ?

$$\begin{aligned} f(x) &= f_0 + f_1x + f_2x^2 + \dots \\ g(x) &= g_0 + g_1x + g_2x^2 + \dots \\ fg &= (f_0 + f_1x + f_2x^2 + \dots)(g_0 + g_1x + g_2x^2 + \dots) \\ &= f_0g_0 + (f_1g_0 + f_0g_1)x + (f_2g_0 + f_1g_1 + f_0g_2)x^2 + \dots \end{aligned}$$

### 8.2.1 Method 1

So, if  $fg = 1$ , we need to simultaneously satisfy

$$\begin{aligned} f_0g_0 = 1 &\Rightarrow f_0 \text{ must be a unit in } \mathbb{R}, \text{ that is } f_0^{-1} \text{ exists so in } R = \mathbb{Z}, \text{ this means } f_0 = \pm 1 \\ f_1g_0 + f_0g_1 &= 0 \end{aligned}$$

$$f_2g_0 + f_1g_1 + f_0g_2 = 0$$

$$\begin{aligned} \text{Then} \quad f_1g_0 = -f_0g_1 & \iff g_1 = -f_0^{-1}f_1g_0 \\ & \implies f_1g_0 + f_0g_1 = 0. \end{aligned}$$

$$\text{Then} \quad f_2g_0 + f_1g_1 + f_0g_2 = 0 \iff g_2 = -f_0^{-1}(f_2g_0 + f_1g_1)$$

$$f_3g_0 + f_2g_1 + f_1g_2 + f_0g_3 = 0 \iff g_3 = -f_0^{-1}(f_3g_0 + f_2g_1 + f_1g_2)$$

So we are able to construct (and actually compute!)  $g(x) = f(x)^{-1}$ .

### 8.2.2 Method 2

**Proof:** Alternative proof: If  $f_0^{-1}$  exist,  $f(x)^{-1}$  exists.

$$\begin{aligned} (1-y)^{-1} &= 1 + y + y^2 \\ f(x) &= f_0(1 + f_0^{-1}f_1x + f_0^{-1}f_2x^2 + \dots) \\ &= f_0(1 - x(-f_0^{-1}f_1 - f_0^{-1}f_2x - f_0^{-1}f_3x^2 + \dots)) \\ &= f_0(1 - xh(x)) \text{ where } h(x) = -f_0^{-1}f_1 - f_0^{-1}f_2x - f_0^{-1}f_3x^2 + \dots \\ f(x)^{-1} &= (1 - xh(x))^{-1} - f_0^{-1} \\ &= (1 + xh(x) + x^2h(x)^2 + x^3h(x)^3 + \dots)f_0^{-1} \end{aligned}$$

which converges since  $|xh(x)|_\mu < 1$ . ■

Example: If  $f(x) = 1 - x - x^{k+1}$ ,

$$f(x)^{-1} = \frac{1}{1 - x(1 + x)^k} = \sum_{k=0}^{\infty} x^\ell (1 + x^k)^\ell$$

$$\begin{aligned} [x^n]f(x)^{-1} &= \sum_{\ell=0}^{\infty} [x^{n-\ell}] (1 + x^k)^\ell \quad \text{Need: } n - \ell \geq 0, k|(n - \ell) \\ &= \sum_{t=0}^{\lfloor \frac{n}{k} \rfloor} [x^{kt}] (1 + x^k)^{n-kt} \quad \text{Put: } kt = n - \ell, 0 \leq kt \leq n \\ &= \sum_{t=0}^{\lfloor \frac{n}{k} \rfloor} \binom{n-kt}{t} \end{aligned}$$

### 8.2.3 Method 3

Under certain circumstances the following is easy-ish to compute. Set  $f(x) = f_0(x)$  and suppose  $f(0) = 1$ .

$$\frac{1}{f_0(x)} = \frac{f_0(-x)}{f_0(x)f_0(-x)}$$

Since  $f_0(x)f_0(-x)$  is even, we can write it as  $f_1(x^2)$ , where perhaps we can compute  $f_1(y)$ .

$$\begin{aligned} \frac{1}{f_0(x)} &= \frac{f_0(-x)}{f_1(x^2)} \\ &= \frac{f_0(-x)f_1(-x^2)}{f_1(x^2)f_1(-x^2)} \\ &= f_0(-x)f_1(-x^2)f_2(-x^4)f_3(-x^8)\cdots \end{aligned}$$

Since  $f_0 = 1 + a_1x + a_2x^2 + \cdots$

$$\begin{aligned} \Rightarrow f_0(x)f_0(-x) &= (1 + a_1x + a_2x^2 + \cdots)(1 + a_1x + a_2x^2 + \cdots) \\ &= 1 + (2a_2 - a_1^2)x^2 + \cdots \\ &= 1 + b_1^2 + b_2x^4 + \cdots \\ \Rightarrow f_k(y) &= 1 + c_1y + c_2y^2 + \cdots \\ f_k(x^{2^k}) &= 1 - c_1x^{2^k} + c_2x^{2^{k+1}} + \cdots \end{aligned}$$

So,  $f_k(x^{2^k}) \rightarrow 1$  as  $k \rightarrow \infty$ .

Note: to compute all coefficients in  $\frac{1}{f(x)}$  up to  $x^N$  requires the product of  $f_k(-x^{2^k})$  up to  $k \geq \log_2 N$  ( $k = \lceil \log_2 N \rceil$ ) giving a product of  $k + 1$  terms. We can do this by using Fourier Transformations.

This is fast precisely when we can compute  $f_k(y)$  efficiently.

**Exercise :**  $f(x) = 1 - x^j$ . What are the  $f_k$ 's? And what does this method give us?

#### 8.2.4 What happens with $f_0 = 1 - x - x^2$ ?

$$\begin{aligned} f_0(-x) &= (1 + x - x^2) \\ f_0(x)f_0(-x) &= (1 + x - x^2)(1 - x - x^2) \\ &= (1 - 3x^2 + x^4) \\ f_1(y) &= 1 - 3y + y^2 \\ f_1(y)f_1(-y) &= (1 - 3y + y^2)(1 + 3y + y^2) \\ &= (1 - 9y^2 + 2y^2 + y^4) \\ &= 1 - 7y^2 + y^4 \\ f_2(z) &= 1 - 7z + z^4 \\ f_k(y) &= (1 - a_k y + y^2) \\ f_k(-y) &= (1 + a_k y + y^2) \\ f_k(y)f_k(-y) &= 1 - (a_k^2 - 2)y^2 + y^4 \\ a_{k+1} &= a_k^2 - 2 \end{aligned}$$

and obtain a recurrence to obtain  $a_k$ 's, and hence a factorization of  $\frac{1}{1-x-x^2}$ .