

# On the number of co-prime-free sets.

by

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**Abstract:** For a variety of arithmetic properties  $P$  (such as the one in the title) we investigate the number of subsets of the positive integers  $\leq x$ , that have that property. In so doing we answer some questions posed by Cameron and Erdős.

## 1. Introduction.

In [CE] Cameron and Erdős investigated subsets of the positive integers  $\leq x$  with certain given properties  $P$ ; in particular, how large such sets can be, and how many there are. The properties  $P$  that they were interested in are monotone decreasing, that is, if  $S$  has property  $P$ , and  $T$  is a subset of  $S$ , then  $T$  has property  $P$ . Thus if  $S$  is a maximal set of positive integers  $\leq x$  with property  $P$  then one knows that there are  $\geq 2^{|S|}$  such sets. In this paper we improve various estimates in [CE] for the number of sets satisfying certain properties  $P$ :

In Theorem 3.5 of [CE], Cameron and Erdős showed that the number of sets of positive integers  $\leq x$ , in which any two elements have a common factor, lies between  $2^{\lfloor x/2 \rfloor}$  and  $x2^{\lfloor x/2 \rfloor}$ . Here we improve this to

**Theorem 1.** *The number of sets of integers  $\leq x$ , with any two elements having a common factor, is*

$$(1.1) \quad 2^{\lfloor x/2 \rfloor} + 2^{\lfloor x/2 \rfloor - N} + O\left(2^{\lfloor x/2 \rfloor - N} \exp\left(-C \frac{x}{\log^2 x \log \log x}\right)\right),$$

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for some absolute constant  $C > 0$ , where  $N$ , which will be defined in the proof, satisfies

$$(1.2) \quad N = (e^{-\gamma} + o(1)) \frac{x}{\log \log x},$$

and  $\gamma$  is the Euler-Mascheroni constant.

In Theorem 3.3 of [CE], Cameron and Erdős showed that the number of sets of positive integers  $\leq x$ , in which any two elements are coprime, lies between  $2^{\pi(x)} e^{(1/2+o(1))\sqrt{x}}$  and  $2^{\pi(x)} e^{(2+o(1))\sqrt{x}}$ . Here we improve this to

**Theorem 2.** *The number of sets of integers  $\leq x$ , with every pair of elements coprime, is*

$$(1.3) \quad 2^{\pi(x)} e^{\sqrt{x}\{1+O(\log \log x/\log x)\}}.$$

In Section 4.3 of [CE], Cameron and Erdős conjectured that there are  $c(s)^{x+o(x)}$  sets of integers  $\leq x$ , with sum of reciprocals bounded by  $s$ , for some positive constant  $c(s)$ . We prove a quantitative form of this conjecture here:

**Theorem 3.** *The number,  $\nu(x)$ , of sets  $\{a_1, a_2, \dots, a_t\}$  of positive integers  $\leq x$ , with  $\sum_i 1/a_i \leq s$ , is*

$$(1.4) \quad c(s)^x e^{O(x^{3/4})}, \quad \text{where } c(s) = \left(1 + e^{-f(s)}\right),$$

and  $f(s)$  is defined by

$$(1.5) \quad s = \int_{f(s)}^{\infty} \frac{du}{u(1+e^u)}.$$

Cameron and Erdős observed that  $c(s) \geq 2^{1-e^{-s}}$ , for all  $s$ . We can provide some rather more accurate estimates based on an analysis of (1.4) and (1.5): As  $s \rightarrow 0$ , we have

$$(1.6) \quad c(s) = 1 + s(\log(1/s) + \log \log(1/s) + O(1));$$

as  $s \rightarrow \infty$ , we have

$$(1.7) \quad c(s) = 2 - Ce^{-2s} + O(e^{-4s}),$$

for some constant  $C \approx .8819384944\dots$

## 2. Sets where any two elements have a common factor.

**Proof of Theorem 1:** Clearly there are  $2^{\lfloor x/2 \rfloor}$  such sets that contain only even numbers.

We now count such sets that contain at least one odd number:

Let  $k$  be the largest integer for which  $q = p_1 p_2 \dots p_k \leq x$  (where  $p_j$  is the  $j$ th smallest odd prime); by the Prime Number Theorem,  $k \sim \log x / \log \log x$ . Define

$$\mathcal{R} = \{n \leq x : n \text{ is divisible by } 2p_j \text{ for some } j \leq k\}.$$

Clearly any set of the form  $S \cup \{q\}$ , where  $S$  is a subset of  $\mathcal{R}$ , has the property that any two elements have a common factor. The number of such subsets  $S$  is  $2^{|\mathcal{R}|}$ , and  $|\mathcal{R}| = \lfloor x/2 \rfloor - N$  where  $N$  is the number of integers  $2m \leq x$  that are coprime to  $q$ . From the combinatorial sieve we obtain  $N \sim \frac{\phi(q)}{q} \frac{x}{2}$ , and then Mertens' Theorem implies (1.2).

Our proof of the rest of Theorem 1 is based on the ideas of Pomerance given in Cameron and Erdős [CE]:

We start by ordering the odd numbers  $\leq x$  as  $q = m_1, m_2, \dots$  so that

$$\frac{\phi(m_1)}{m_1} \leq \frac{\phi(m_2)}{m_2} \leq \frac{\phi(m_3)}{m_3} \leq \dots$$

Define  $f_i(x)$  to be the number of sets of integers  $\leq x$ , that contain  $m_i$  but not  $m_{i+1}, m_{i+2}, \dots$ , and for which every pair of integers in the set have a common factor.

If  $\frac{\phi(m_i)}{m_i} \geq 2/3$  then

$$\#\{n \leq x : (n, m_i) > 1\} \leq \sum_{p|m_i} \frac{x}{p} \leq x \sum_{p|m_i} \log \left( \frac{p}{p-1} \right) = x \log \left( \frac{m_i}{\phi(m_i)} \right) \leq x \log (3/2),$$

so that  $f_i(x) \leq 2^{x \log (3/2)}$ , which is part of the error term in (1.1).

Define  $A_m(x)$  to be the number of even integers  $\leq x$ , that are coprime to  $m$ , and  $c_i = A_{m_i}(x)$ . Clearly  $f_i(x) \leq 2^{i-1+[x/2]-c_i}$ , and so, in order to complete the proof of Theorem 1, we must show that

$$(2.1) \quad c_i(x) - i - N \gg \frac{x}{\log^2 x \log \log x},$$

for all  $i \geq 2$  for which  $\frac{\phi(m_i)}{m_i} \leq 2/3$ .

Suppose that  $m = pr \leq m' = p'r \leq x$  are odd, squarefree numbers, with  $p < x^{1/3}$ . Noting that  $|A_r(x)| = |A_m(x)| + |A_r(x/p)|$ , we see that

$$(2.2) \quad |A_{m'}(x)| - |A_m(x)| = |A_r(x/p) \setminus A_r(x/p')| \gg \left(\frac{1}{p} - \frac{1}{p'}\right) \frac{x}{\log \log x} \gg \frac{x}{p^2 \log \log x}.$$

Thus, given any squarefree integer  $r \leq x$ ,  $r \neq q$ , we form a sequence  $r = r_0, r_1, \dots, r_j = q$ , as follows: If  $r_i = p_1 p_2 \dots p_h$  for some  $h < k$ , then then let  $r_{i+1} = r_i p_{h+1}$ . Otherwise we construct  $r_{i+1}$  by dividing  $r_i$  by its largest prime factor, and multiplying it by the smallest odd prime that does not yet divide it.

Now as  $|A_q(x)| = N$ , and as the prime factors of  $q$  are all  $\ll \log x$ , we find that, if  $i \geq 2$  then  $c_i - N \gg x/\log^2 x \log \log x$  by (2.2); which implies (2.1) for  $i \ll x/\log^2 x \log \log x$ .

So we are left with those  $i$  for which  $\frac{\phi(m_i)}{m_i} \leq 2/3$  and  $i \gg x/\log^2 x \log \log x$ . We first deal with those  $i \leq 100x/\log \log x$ :

If  $\mathcal{A}$  is any set of  $< i$  odd integers  $\leq x$  then

$$\frac{\phi(m_i)}{m_i} \geq \min_{\substack{n \leq x, \\ n \notin \mathcal{A}}} \frac{\phi(n)}{n}.$$

So we choose  $\mathcal{A}$  to be the set of those  $[x/\log^3 x]$  odd integers  $\leq x$ , with the most distinct prime factors.

Hardy and Ramanujan [HR] showed that there exists an absolute constant  $c$  such that the number of integers  $\leq x$  with exactly  $k$  distinct prime factors is

$$\ll \frac{x}{\log x} \frac{(\log \log x + c)^{k-1}}{(k-1)!};$$

therefore the number of integers  $\leq x$  with at least  $10\log\log x$  distinct prime factors is  $\ll x/\log^{10}x$ . Thus, if  $n \notin \mathcal{A}$  then  $n$  has  $\leq 10\log\log x$  distinct prime factors, and so

$$\frac{\phi(n)}{n} \geq \prod_{p \leq (\log\log x)^2} \left(1 - \frac{1}{p}\right) \gg \frac{1}{\log\log\log x}$$

by Mertens' Theorem. Therefore, by the combinatorial sieve,

$$c_i \gg \frac{\phi(m_i)}{m_i} x \gg \frac{x}{\log\log\log x},$$

which implies (2.1) in this range of  $i$ .

Finally we come to those  $i$  in the range  $100x/\log\log x \leq i \leq x/2$ , with  $\frac{\phi(m_i)}{m_i} \leq 2/3$ :

An immediate consequence of Proposition 4 of [PS] is that

$$i < \frac{3}{20} x \frac{\phi(m_i)}{m_i} / \left(1 - \frac{\phi(m_i)}{m_i}\right) \leq \frac{9}{20} x \frac{\phi(m_i)}{m_i}.$$

On the other hand, by the combinatorial sieve,

$$c_i \geq \left\{ \frac{1}{2} + o(1) \right\} x \frac{\phi(m_i)}{m_i} > \left\{ \frac{10}{9} + o(1) \right\} i;$$

and thus

$$c_i - i \geq \frac{i}{10} \geq \frac{10x}{\log\log x},$$

and so (2.1) is satisfied.

### 3. Sets where any two elements are coprime.

**Proof of Theorem 2:** Let  $t = \pi(\sqrt{x})$ . For the upper bound, note that the number of composite elements of each set is  $\leq t$ , as these elements must all have distinct prime factors  $\leq x$ . All other elements are prime, and so the number of such sets is

$$\leq 2^{\pi(x)} \sum_{i=0}^t \binom{[x]}{i} \ll 2^{\pi(x)} \frac{x^t}{t!} \ll 2^{\pi(x)} \left(\frac{ex}{t}\right)^t,$$

which gives (1.3) by the Prime Number Theorem. (The proof here is the same as in Theorem 3.3 of [CE], except that they made a computational error in the final step.)

To obtain the lower bound, we shall construct (1.3) such sets. Let

$$k = t \left( 1 - \frac{\log \log x}{\log x} \right)$$

and let

$$\sqrt{x} < q_1 < q_2 < \dots < q_k$$

be the  $k$  smallest primes larger than  $\sqrt{x}$ . Note that, using the Prime Number Theorem in the form  $\pi(x) = \frac{x}{\log x}(1 + O(1/\log x))$  we have

$$(3.1) \quad q_j = x^{1/2}(1 + j/t + O(1/\log x))$$

and so  $q_k < 2\sqrt{x}$ .

We construct our sets as follows:

Each prime in the interval  $(2\sqrt{x}, x]$  is in our set or not as desired, giving  $2^{\pi(x) - \pi(2\sqrt{x})}$  different options.

We may put any number of the form  $p_k q_k$  in the set, where  $p_k$  is a prime less than  $x/q_k$  (giving  $\pi(x/q_k)$  choices).

Then any  $p_{k-1} q_{k-1}$  where  $p_{k-1}$  is a prime  $\leq x/q_{k-1}$  (giving  $\pi(x/q_{k-1}) - 1$  choices).

We continue in this fashion, taking, in general, any  $p_{k-j} q_{k-j}$ , where  $p_{k-j}$  is a prime  $\leq x/q_{k-j}$ , not already used as some  $p_{k-i}$ , (giving us  $\pi(x/q_{k-j}) - j$  choices), for  $j = 0, 1, \dots, k-1$ .

Thus the number of different sets constructed is

$$2^{\pi(x) + O(\sqrt{x}/\log x)} \prod_{i=0}^k \left\{ \pi \left( \frac{x}{q_i} \right) - (k-i) \right\}.$$

Now, by (3.1),

$$\begin{aligned} \pi \left( \frac{x}{q_i} \right) - (k-i) &= \pi \left( \frac{x^{1/2}}{1 + i/t} \left( 1 + O \left( \frac{1}{\log x} \right) \right) \right) - (k-i) \\ &= t \left\{ \frac{1}{1 + i/t} + O \left( \frac{1}{\log x} \right) - 1 + \frac{\log \log x}{\log x} + \frac{i}{t} \right\}, \end{aligned}$$

using the Prime Number Theorem again,

$$= t \left( \frac{\log \log x}{\log x} + \frac{(i/t)^2}{1+i/t} + O\left(\frac{1}{\log x}\right) \right) \geq \frac{t}{\log x}.$$

Therefore, the number of sets is at least

$$2^{\pi(x)+O(\sqrt{x}/\log x)} \left( \frac{t}{\log x} \right)^k,$$

which gives (1.3).

**Remark:** With some care it is possible to replace the  $\log \log x$  in (1.3) with  $(\log \log \log x)^2$ , but we are currently unable to do better.

#### 4. Sets whose sum of reciprocals is bounded.

**Proof of Theorem 3:** Let  $y = [x^{1/4}]$ ,  $z = [x^{1/2}]$ , and  $x_j = jx/z$  for  $y \leq j \leq z$ . A given set of positive integers  $\leq x$ , whose sum of reciprocals is bounded by  $s$ , has, say,  $b_j$  integers in the interval  $(x_j, x_{j+1}]$  for  $y \leq j < z$  (with  $0 \leq b_j \leq x/z + 1$ ), and so satisfies

$$(4.1) \quad \sum_{j=y}^{z-1} \frac{b_j}{x_{j+1}} \leq s.$$

So, if the  $b_j$  are fixed with these values then the number of sets of integers  $\leq x$ , with precisely  $b_j$  integers from the interval  $(x_j, x_{j+1}]$ , is

$$(4.2) \quad \leq 2^{xy/z} \prod_{y \leq j < z} \binom{[x/z] + 1}{b_j}.$$

So define  $\alpha_j := b_j / ([x/z] + 1)$  for each  $j$ . Clearly  $0 \leq \alpha_j \leq 1$  for each  $j$ , and, by (4.1), they must satisfy

$$(4.3) \quad \sum_{j=y}^{z-1} \frac{\alpha_j}{j} \leq s \left( 1 + O\left(\frac{1}{y}\right) \right).$$

Moreover, by Stirling's formula,

$$\binom{[x/z] + 1}{b_j} \ll \exp\left(O(\log x) - \frac{x}{z}(\alpha_j \log \alpha_j + (1 - \alpha_j) \log(1 - \alpha_j))\right).$$

Noting that there are no more than  $(x/z + 1)^z$  choices for the  $b_j$ , we thus see that

$$\nu(x) \ll \exp\left(O(x^{3/4}) - \frac{x}{z} \min_{\substack{0 \leq \alpha_j \leq 1, y \leq j < z \\ (4.3) \text{ holds}}} \sum_{y \leq j < z} (\alpha_j \log \alpha_j + (1 - \alpha_j) \log(1 - \alpha_j))\right).$$

By the method of Lagrange multipliers we find that the minimum occurs when each  $\alpha_j = 1/(1 + e^{A/j})$  for some constant  $A > 0$ .

Now

$$\begin{aligned} \sum_{j=y}^{z-1} \frac{\alpha_j}{j} &= \int_y^z \frac{dt}{t(1 + e^{A/t})} + O\left(\frac{1}{ye^{A/y}} + \frac{1}{ze^{A/z}}\right) \\ &= \int_{A/z}^{\infty} \frac{du}{u(1 + e^u)} + O\left(\frac{1}{ye^{A/y}} + \frac{1}{(A/y)e^{A/y}} + \frac{1}{ze^{A/z}}\right). \end{aligned}$$

To obtain equality in (4.3), we need to select  $A = x^{1/2}f(s) + O(x^{1/4})$ . Thus

$$\begin{aligned} - \sum_{y \leq j < z} \alpha_j \log \alpha_j + (1 - \alpha_j) \log(1 - \alpha_j) &= \sum_{y \leq j < z} \log\left(1 + e^{-A/j}\right) + \frac{A}{j} \frac{1}{(1 + e^{A/j})} \\ &= \int_y^z \left( \left(1 + e^{-A/u}\right) + \frac{A}{u(1 + e^{A/u})} \right) du + O(1). \end{aligned}$$

By the substitution  $t = A/u$ , and the fact that

$$\frac{d}{dt} \frac{\log(1 + e^{-t})}{t} = \frac{1}{t^2} \left( \log(1 + e^{-t}) + \frac{t}{1 + e^t} \right),$$

this last line is

$$z \log(1 + e^{-A/z}) + O(1),$$

and so (1.4) is an upper bound for  $\nu(x)$ .



To obtain a lower bound for  $\nu(x)$  note that if we select integers  $b_j$  such that

$$(4.4) \quad \sum_{j=y}^{z-1} \frac{b_j}{x_j} \leq s,$$

then the sum of the reciprocals of any set consisting of  $b_j$  integers from the interval  $(x_j, x_{j+1})$ , for each  $y \leq j < z$ , is  $\leq s$ . Clearly the number of such sets is

$$(4.5) \quad \geq \prod_{y \leq j < z} \binom{[x/z]}{b_j}.$$

Now the idea is to select each  $b_j = x/z(1 + e^{A/j}) + O(1)$ , for some constant  $A > 0$ , so that (4.4) is satisfied; this can be done by the choice  $A = x^{1/2}f(s) + O(x^{1/4})$  (the proof being almost identical to that for the upper bound). Now, when we estimate (4.5) with this value for  $A$ , we proceed as in the upper bound, and show that (4.5) is at least (1.4).

**Remarks:** Cameron and Erdős observed that any set of integers taken from  $[x/e^s, x]$  has sum of reciprocals  $\leq s$ , and thus  $c(s) \geq 2^{1-e^{-s}}$ . We will derive (1.6) and (1.7):

If  $x$  is ‘large’ then

$$\int_x^\infty \frac{du}{u(1+e^u)} = \frac{1}{xe^x} \left\{ 1 + O\left(\frac{1}{x}\right) \right\},$$

which implies (1.6). On the other hand, (1.5) gives that, for  $f(s) < 1$  (that is,  $s$  ‘large’),

$$(4.6) \quad s = C_1 + \frac{1}{2} \log(1/f(s)) - \int_0^{f(s)} \left( \frac{1}{1+e^u} - \frac{1}{2} \right) \frac{du}{u},$$

where

$$C_1 = \int_1^\infty \frac{du}{u(1+e^u)} + \int_0^1 \left( \frac{1}{1+e^u} - \frac{1}{2} \right) \frac{du}{u}.$$

Therefore  $f(s) \asymp e^{-2s}$ , and so the last term in (4.6) is  $\ll f(s) \ll e^{-2s}$ . Thus  $f(s) = Ce^{-2s} + O(e^{-4s})$ , where  $C = e^{2C_1}$  (which can be computed explicitly), which implies (1.7).

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