

MTHSC 412 SECTION 2.5 –CONGRUENCE OF INTEGERS

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CONGRUENCE MODULO n IS AN EQUIVALENCE RELATION ON \mathbb{Z}

DEFINITION

Let $n > 1$ be an integer. For $x, y \in \mathbb{Z}$, we say that x is congruent to y modulo n and write $x \equiv y \pmod{n}$ if $n \mid (x - y)$.

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If $n > 1$ is an integer then $\equiv \pmod{n}$ is an equivalence relation on \mathbb{Z} .

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Then $(x - y) = nk$ and $(y - z) = nm$ for some $k, m \in \mathbb{Z}$.

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So, $(x - z) = (x - y) + (y - z) = n(k + m)$ and $n|(x - z)$.

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Thus $x \equiv z \pmod{n}$ and $\equiv \pmod{n}$ is transitive.

Since $\equiv \pmod{n}$ is reflexive, symmetric and transitive, it is an equivalence relation on \mathbb{Z} . □

CONGRUENCE AND REMAINDERS

FACT

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Now note that $n \mid (x - y)$ if and only if $n \mid (r_1 - r_2)$.

Finally since $-n < (r_1 - r_2) < n$, $n \mid (r_1 - r_2)$ if and only if $r_1 = r_2$. □

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Thus each integer is in one of the congruence classes:

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The fact that these are distinct follows from our last fact. □

ADDITION AND MULTIPLICATION PROPERTIES

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If $a \equiv b \pmod{n}$ and $x \in \mathbb{Z}$ then

$$a + x \equiv b + x \pmod{n} \quad \text{and} \quad ax \equiv bx \pmod{n}.$$

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PROOF.

Suppose that $a \equiv b \pmod{n}$. Then $(a - b) = nk$ for some $k \in \mathbb{Z}$. Thus $(a + x) - (b + x) = a - b = nk$ and $ax - bx = x(a - b) = xnk$ and the result follows. \square

THEOREM

Suppose that $a \equiv b \pmod{n}$ and $c \equiv d \pmod{n}$. Then

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Thus, $ac \equiv bc \equiv bd \pmod{n}$.

The proof of the other congruence is similar and is left as an exercise. □

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$$\begin{aligned} ax \equiv ay \pmod{n} &\Rightarrow n \mid (ax - ay) \\ &\Rightarrow n \mid a(x - y) \quad \text{and} \quad (a, n) = 1 \end{aligned}$$

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 \end{aligned}$$



THEOREM

If $(a, n) = 1$, the congruence $ax \equiv b \pmod{n}$ has a solution $x \in \mathbb{Z}$ and the solution is unique modulo n , which means that any two such solutions are congruent modulo n .

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Uniqueness modulo n : Suppose that $x, y \in \mathbb{Z}$ are both solutions.

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Uniqueness modulo n : Suppose that $x, y \in \mathbb{Z}$ are both solutions. Then $ax \equiv b \equiv ay \pmod{n}$ and $(a, n) = 1$.

By the cancellation law, it follows that $x \equiv y \pmod{n}$. □

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COROLLARY

If n_1, n_2, \dots, n_k is a set of pairwise coprime integers and if $n_i|c$ for $1 \leq i \leq k$, then $(n_1 n_2 \dots n_k)|c$.

CHINESE REMAINDER THEOREM

THEOREM

Let n_1, n_2, \dots, n_k be pairwise coprime integers. Let $a_1, \dots, a_k \in \mathbb{Z}$. There is $x \in \mathbb{Z}$ satisfying the system of congruences

$$\begin{cases} x \equiv a_1 \pmod{n_1} \\ x \equiv a_2 \pmod{n_2} \\ \vdots \\ x \equiv a_k \pmod{n_k}. \end{cases}$$

Furthermore, the solution is unique modulo $(n_1 n_2 \dots n_k)$.

PROOF.

Existence: Let $N = (n_1 n_2 \dots n_k)$ and let
 $N_i = \frac{N}{n_i} = n_1 n_2 \dots n_{i-1} n_{i+1} \dots n_k.$

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Let M_i be a solution to $N_i x \equiv 1 \pmod{n_i}$.

$$\text{Then } N_i M_i \equiv \begin{cases} 0 \pmod{n_j} & \text{if } i \neq j, \\ 1 \pmod{n_i} \end{cases}.$$

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Now let $x = \sum_{i=1}^k a_i N_i M_i$.

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$$N_i = \frac{N}{n_i} = n_1 n_2 \dots n_{i-1} n_{i+1} \dots n_k.$$

Then $(N_i, n_i) = 1$ because the n_i are pairwise coprime.

Let M_i be a solution to $N_i x \equiv 1 \pmod{n_i}$.

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Thus $x \equiv y \pmod{N}$ as desired. □