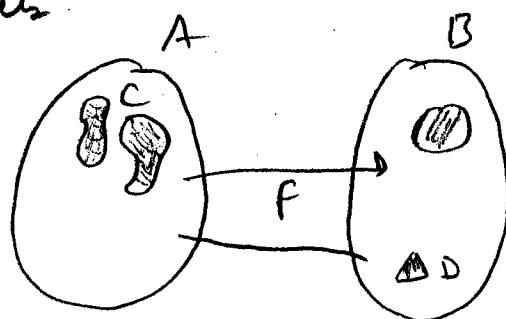


Lecture 7&8 Countable & uncountable sets

How to count:

domain codomain

- Recall: $f: A \rightarrow B$
"maps" $x \mapsto f(x)$.



If $C \subseteq A$, $D \subseteq B$, define $f(C) = \{f(x) : x \in C\}$, the image of C .

and $f^{-1}(D) = \{x \in A : f(x) \in D\}$, the inverse image, or preimage of D .

Note that f^{-1} need not be a function.

Def: • $f(A)$ is the range of f . When $f(A) = B$, say f

(i) onto (or surjective). Write $f: A \rightarrow B$

write $f: A \hookrightarrow B$.

• When $f(x) = f(y)$ implies $x = y$, say f is 1-1 (or injective)

• When f is 1-1 & onto, call f a bijection, and we say

$A \& B$ are in 1-1 correspondence. Write $A \sim B$.

Elementary counting: Use $A = J_n = \{1, 2, \dots, n\}$.

Ex: $\{1, a, 88, \emptyset\} = \boxed{A}$. Say $|A| = 4$.

$\begin{matrix} \uparrow & \uparrow & \uparrow & \uparrow & \uparrow \\ 3 & 1 & 4 & 2 & J_4 \end{matrix}$

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Def: Call A finite if $A \sim J_n$ for some n . [Call $J_0 = \emptyset$].

Otherwise, call A infinite.

Def: Call A countable if $A \sim \mathbb{N}$. [Sometimes, we include finite sets too.]

Ex: \mathbb{N} is countable: use $f: \mathbb{N} \rightarrow \mathbb{N}$, $f(x) = x$.

• A sequence x_1, x_2, x_3, \dots of distinct terms is countable:

use $f: \mathbb{N} \rightarrow \{x_n\}$ $f(n) = x_n$.

* Moral: A set that can be "listed" in a sequence is countable!

Ex: $\{2, 3, 4, 5, \dots\}$ is countable, use $f(n) = n+1$

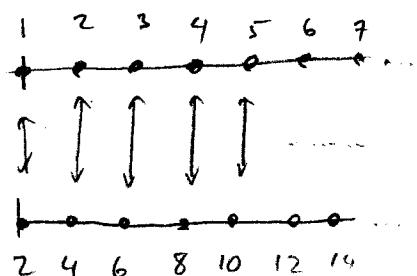
• $\{1, 2, 3, \dots, k-1, k+1, k+2, \dots\}$ is countable: use $f(n) = \begin{cases} n & n < k \\ n+1 & n \geq k \end{cases}$

Theorem: \mathbb{N} is infinite.

Proof (Sketch) Induction. I'll skip this.

Ex: $\mathbb{N} \sim 2\mathbb{N} = \{2, 4, 6, 8, 10, \dots\}$. Use $f(n) = 2n$.

$\mathbb{N}, 2\mathbb{N}$ have the same cardinality



Ex: $\mathbb{Z} = \{\dots, -3, -2, -1, 0, 1, 2, 3, \dots\}$ is countable.

... 7 5 3 1 2 4 6 ...

3

Thm: Every infinite subset E of a countable set A is countable.

Proof: (sketch). $A = \{x_1, x_2, x_3, x_4, x_5, \dots\}$

$$\text{let } n_1 = \inf\{i : x_i \in E\}$$

$$n_2 = \inf\{i : x_i \in E, i > n_1\}$$

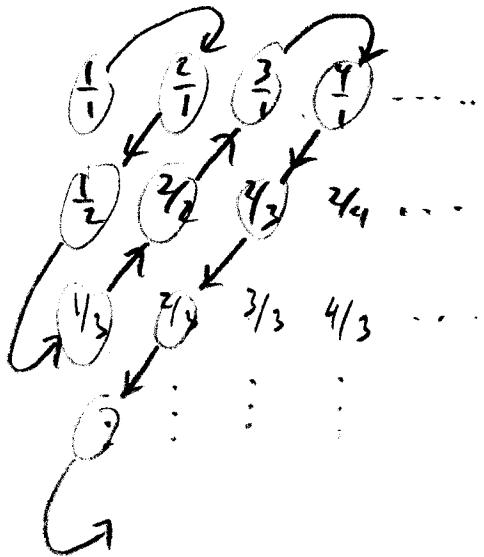
⋮

$$n_k = \inf\{i : x_i \in E, i > n_{k-1}\}. \quad \text{Then } E = \{x_{i_1}, x_{i_2}, x_{i_3}, \dots\}.$$

□

Thm: \mathbb{Q} is countable.

Proof:



Remark: This shows that ordered pairs, (m, n) , $m, n \in \mathbb{Z}$ are countable.

- \mathbb{Q}^+ is a subset of this
- It's easy to extend this to all of \mathbb{Q} .

Thm: If A is countable, then $A \times A$ is countable.

Theorem: \mathbb{R} is not countable.

Proof: (by contradiction). Suppose \exists bijection $f: \mathbb{N} \rightarrow \mathbb{R}$.

Consider the list $f(1), f(2), f(3), \dots$

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$f: \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}^{\omega} \rightarrow \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}^{\omega}$ [Disallow infinite strings of 9]

$f(1) \mapsto 0.1234567890\dots$

$f(2) \mapsto 0.3234170092\dots$

$f(3) \mapsto 0.3141592653\dots$

$f(4) \mapsto 0.9871243882$

$f(5) \mapsto 0.8767234100$ increase each digit by $\begin{cases} 1 & \text{if it's } 0-7 \\ -1 & \text{if it's } 8-9 \end{cases}$.

Construct a new number $x^* = 0.23523\dots$

Note that $x^* \neq f(1)$ (differs in 1st digit)

$x^* \neq f(2)$ (differs in 2nd digit)

$x^* \neq f(3)$ (differs in 3rd digit).

so x^* cannot be $f(n)$ for any n , contradicting the assumption that f is surjective. \square

Sets of different cardinalities: $\mathbb{N}, \mathbb{Z}, \dots, \mathbb{J}_n, \dots, \mathbb{N}, \mathbb{R}, ? \dots$

Remark: The question of if there is a set A with cardinality $|\mathbb{N}| < |A| < |\mathbb{R}|$ is undecidable, i.e., it is neither true nor false! This is the continuum hypothesis.

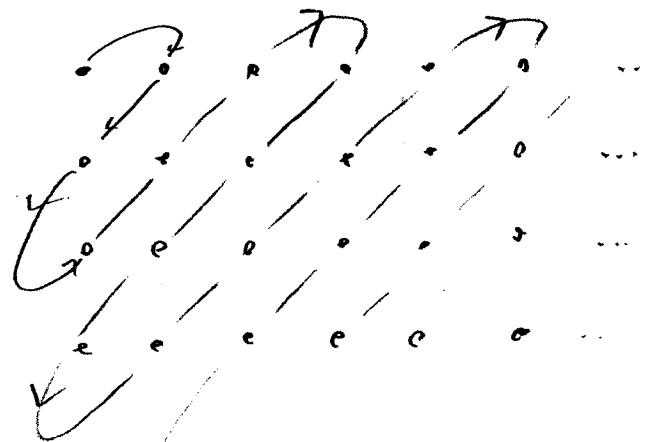
Theorem: The countable union of countable sets is countable.

Proof: Say each A_1, A_2, A_3, \dots is countable

Then $A_1 = \{a_{11}, a_{12}, a_{13}, a_{14}, a_{15}, \dots\}$

$A_2 = \{a_{21}, a_{22}, a_{23}, a_{24}, a_{25}, \dots\}$

Apply Cantor's diagonal argument:



Remark: We can replace "countable" with "at most countable". D

Notation: Use $\bigcup_{\alpha \in J} A_\alpha$ for possibly uncountable collection J .

Ex: • The set of computer programs is countable!

• The set of all possible outputs to computer programs is thus (at most) countable.

But the set \mathbb{R} is uncountable!

* Thus, there are real numbers that are not "computable"! (In fact, most!)

i.e., input: $k \in \mathbb{N}$, output: k decimal places.

• Given a set A , the powerset 2^A is the set of all subsets of A .

Example: $A = \{1, a, 880\}$ $2^A = \{\emptyset, \{1\}, \{a\}, \{880\}, \{1, a\}, \{1, 880\}, \dots\}$

Encode B as $(0|1)$ $1 = \text{yes}$, $a = \text{no}$, $880 = \text{yes}$ ↑ call this B .

There is a bijection $2^A \rightarrow$ binary strings of length $2^{|A|}$.

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Cantor's Theorem (diagonalization argument, 1891):

For any set A , we have $A \not\sim 2^A$

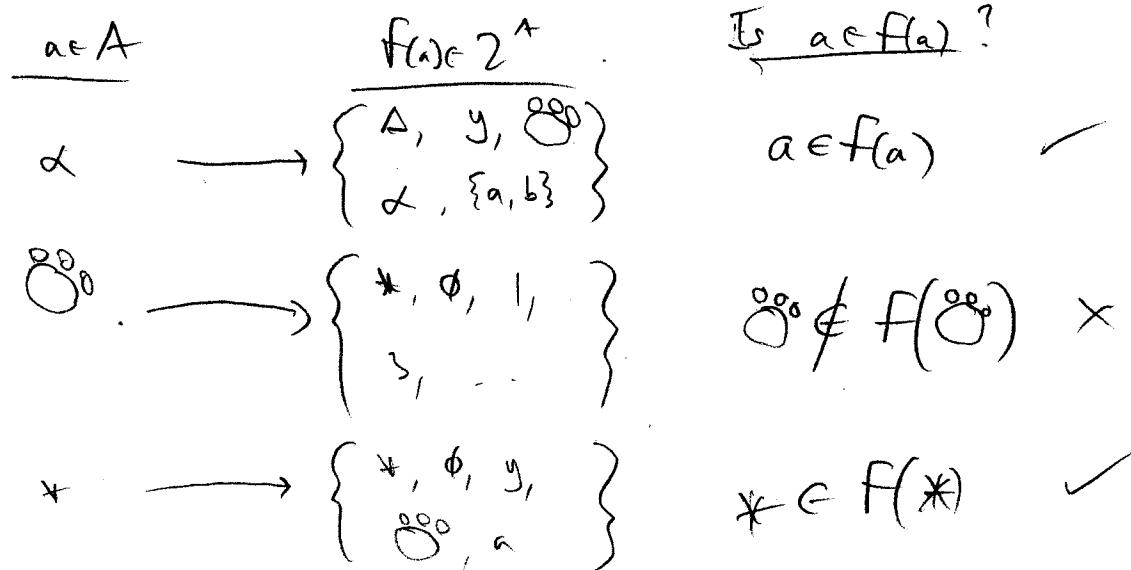
Proof: (by contradiction).

Suppose \exists bijection $f: A \leftrightarrow 2^A$

$$a \mapsto f(a) \subset A.$$

$$\text{e.g., } a \mapsto \{1, \aleph_0, a\}$$

Idea:



Want subset $B \in 2^A$ not in image of f .

$$\text{Let } B = \{a : a \notin f(a)\}.$$

So, if $B = f(x)$ for some $x \in A$, consider x . ($x \xrightarrow{f} B$)

* Is $x \in B$? No, because then $x \notin f(x) = B$, contradiction

Then $x \notin B$, then $x \in f(x) \setminus B$, so $x \in B$, contradiction

So $B \neq f(x)$ for any $x \in A$, thus such a bijection can't exist. □

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Consider the set $F_A = \{f: A \rightarrow \{0,1\}\}$, where A is any set.

Claims $2^A \sim F_A$.

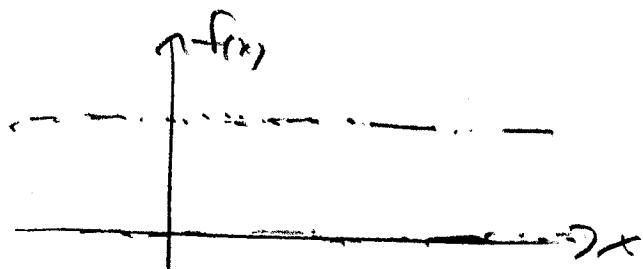
Proof (sketch). There is a Bijection

$$2^A \rightarrow F_A$$

$$S \mapsto f(x) = \begin{cases} 1 & x \in S \\ 0 & x \notin S. \end{cases}$$

Example: If $A = \mathbb{R}$, then F_A is the set of all real-valued functions $\mathbb{R} \rightarrow \{0,1\}$

There are $2^{\mathbb{R}}$ such functions!



Questions:

- How big is the set $\{f: \mathbb{R} \rightarrow \mathbb{R}\}$?

- How big is the set $\{f: \mathbb{R} \rightarrow \mathbb{R}, f \text{ is continuous}\}$?

We'll discuss these later...