## Lecture 5.1: Symmetric cryptographic ciphers

#### Matthew Macauley

Department of Mathematical Sciences Clemson University http://www.math.clemson.edu/~macaule/

Math 4190, Discrete Mathematical Structures

In this lecture, we'll see how to send encoded messages, which will be numbers.

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

A B C D E F G H I J K L M N O P Q R S T U V W X Y Z 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

A B C D E F G H I J K L M N O P Q R S T U V W X Y Z 0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25

In base 26, the word  ${f CLEMSON}$  can be encoded as 2 11 4 12 18 14 13.

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

In base 26, the word **CLEMSON** can be encoded as 2 11 4 12 18 14 13.

We can convert this to decimal (base 10):

$$2 \cdot 26^0 + 11 \cdot 26^1 + 4 \cdot 26^2 + 12 \cdot 26^3 + 18 \cdot 26^4 + 14 \cdot 26^5 + 13 \cdot 26^6 = 4190683824.$$

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

In base 26, the word **CLEMSON** can be encoded as 2 11 4 12 18 14 13.

We can convert this to decimal (base 10):

$$2 \cdot 26^0 + 11 \cdot 26^1 + 4 \cdot 26^2 + 12 \cdot 26^3 + 18 \cdot 26^4 + 14 \cdot 26^5 + 13 \cdot 26^6 = 4190683824.$$

To reverse this process, recursively divide 26 into the number.

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

In base 26, the word **CLEMSON** can be encoded as 2 11 4 12 18 14 13.

We can convert this to decimal (base 10):

$$2 \cdot 26^0 + 11 \cdot 26^1 + 4 \cdot 26^2 + 12 \cdot 26^3 + 18 \cdot 26^4 + 14 \cdot 26^5 + 13 \cdot 26^6 = 4190683824.$$

$$221707947 = 8527228 \cdot 26 + 19$$

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

In base 26, the word **CLEMSON** can be encoded as 2 11 4 12 18 14 13.

We can convert this to decimal (base 10):

$$2 \cdot 26^0 + 11 \cdot 26^1 + 4 \cdot 26^2 + 12 \cdot 26^3 + 18 \cdot 26^4 + 14 \cdot 26^5 + 13 \cdot 26^6 = 4190683824.$$

$$221707947 = 8527228 \cdot 26 + 19$$
 **T**

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

In base 26, the word **CLEMSON** can be encoded as 2 11 4 12 18 14 13.

We can convert this to decimal (base 10):

$$2 \cdot 26^0 + 11 \cdot 26^1 + 4 \cdot 26^2 + 12 \cdot 26^3 + 18 \cdot 26^4 + 14 \cdot 26^5 + 13 \cdot 26^6 = 4190683824.$$

$$221707947 = 8527228 \cdot 26 + 19$$
 **T**  $8527228 = 327970 \cdot 26 + 8$ 

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

In base 26, the word **CLEMSON** can be encoded as 2 11 4 12 18 14 13.

We can convert this to decimal (base 10):

$$2 \cdot 26^0 + 11 \cdot 26^1 + 4 \cdot 26^2 + 12 \cdot 26^3 + 18 \cdot 26^4 + 14 \cdot 26^5 + 13 \cdot 26^6 = 4190683824.$$

$$221707947 = 8527228 \cdot 26 + 19$$
 **T**  $8527228 = 327970 \cdot 26 + 8$  **I**

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

In base 26, the word  ${f CLEMSON}$  can be encoded as 2 11 4 12 18 14 13.

We can convert this to decimal (base 10):

$$2 \cdot 26^0 + 11 \cdot 26^1 + 4 \cdot 26^2 + 12 \cdot 26^3 + 18 \cdot 26^4 + 14 \cdot 26^5 + 13 \cdot 26^6 = 4190683824.$$

$$221707947 = 8527228 \cdot 26 + 19$$
 **T**
 $8527228 = 327970 \cdot 26 + 8$  **I**
 $327970 = 12614 \cdot 26 + 6$ 

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

In base 26, the word  ${f CLEMSON}$  can be encoded as 2 11 4 12 18 14 13.

We can convert this to decimal (base 10):

$$2 \cdot 26^0 + 11 \cdot 26^1 + 4 \cdot 26^2 + 12 \cdot 26^3 + 18 \cdot 26^4 + 14 \cdot 26^5 + 13 \cdot 26^6 = 4190683824.$$

$$221707947 = 8527228 \cdot 26 + 19$$
 **T**
 $8527228 = 327970 \cdot 26 + 8$  **I**
 $327970 = 12614 \cdot 26 + 6$  **G**

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

In base 26, the word **CLEMSON** can be encoded as 2 11 4 12 18 14 13.

We can convert this to decimal (base 10):

$$2 \cdot 26^0 + 11 \cdot 26^1 + 4 \cdot 26^2 + 12 \cdot 26^3 + 18 \cdot 26^4 + 14 \cdot 26^5 + 13 \cdot 26^6 = 4190683824.$$

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

In base 26, the word **CLEMSON** can be encoded as 2 11 4 12 18 14 13.

We can convert this to decimal (base 10):

$$2 \cdot 26^0 + 11 \cdot 26^1 + 4 \cdot 26^2 + 12 \cdot 26^3 + 18 \cdot 26^4 + 14 \cdot 26^5 + 13 \cdot 26^6 = 4190683824.$$

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

In base 26, the word **CLEMSON** can be encoded as 2 11 4 12 18 14 13.

We can convert this to decimal (base 10):

$$2 \cdot 26^0 + 11 \cdot 26^1 + 4 \cdot 26^2 + 12 \cdot 26^3 + 18 \cdot 26^4 + 14 \cdot 26^5 + 13 \cdot 26^6 = 4190683824.$$

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

In base 26, the word **CLEMSON** can be encoded as 2 11 4 12 18 14 13.

We can convert this to decimal (base 10):

$$2 \cdot 26^0 + 11 \cdot 26^1 + 4 \cdot 26^2 + 12 \cdot 26^3 + 18 \cdot 26^4 + 14 \cdot 26^5 + 13 \cdot 26^6 = 4190683824.$$

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

In base 26, the word **CLEMSON** can be encoded as 2 11 4 12 18 14 13.

We can convert this to decimal (base 10):

$$2 \cdot 26^0 + 11 \cdot 26^1 + 4 \cdot 26^2 + 12 \cdot 26^3 + 18 \cdot 26^4 + 14 \cdot 26^5 + 13 \cdot 26^6 = 4190683824.$$

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

In base 26, the word **CLEMSON** can be encoded as 2 11 4 12 18 14 13.

We can convert this to decimal (base 10):

$$2 \cdot 26^0 + 11 \cdot 26^1 + 4 \cdot 26^2 + 12 \cdot 26^3 + 18 \cdot 26^4 + 14 \cdot 26^5 + 13 \cdot 26^6 = 4190683824.$$

In this lecture, we'll see how to send encoded messages, which will be numbers.

We can encode any word as a number in base-26:

In base 26, the word  ${f CLEMSON}$  can be encoded as 2 11 4 12 18 14 13.

We can convert this to decimal (base 10):

$$2 \cdot 26^0 + 11 \cdot 26^1 + 4 \cdot 26^2 + 12 \cdot 26^3 + 18 \cdot 26^4 + 14 \cdot 26^5 + 13 \cdot 26^6 = 4190683824.$$

To reverse this process, recursively divide 26 into the number. Let's try this with 221707947:

Now, suppose that we wanted to send this as a secret message...

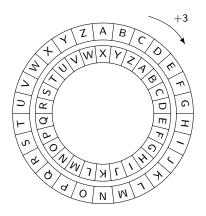
Though he wasn't the first, Julius Caesar (100 B.C-44 B.C) used an encryption device called a cipher in his private correspondences.

Though he wasn't the first, Julius Caesar (100 B.C–44 B.C) used an encryption device called a cipher in his private correspondences.

An encrypted message would looks something like this: RZ WKDWV VKDUS

Though he wasn't the first, Julius Caesar (100 B.C–44 B.C) used an encryption device called a cipher in his private correspondences.

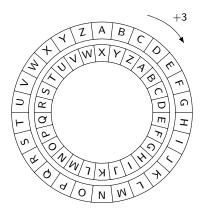
An encrypted message would looks something like this: RZ WKDWV VKDUS



Decrypted message:

Though he wasn't the first, Julius Caesar (100 B.C–44 B.C) used an encryption device called a cipher in his private correspondences.

An encrypted message would looks something like this: RZ WKDWV VKDUS



Decrypted message: OW THATS SHARP

The Caesar cipher is defined by the following:

 $\blacksquare$  key,  $k \in \mathbb{N}$ ,

The Caesar cipher is defined by the following:

- key,  $k \in \mathbb{N}$ ,
- encryption function, e(x),

The Caesar cipher is defined by the following:

- key,  $k \in \mathbb{N}$ ,
- encryption function, e(x),
- decryption function, d(y),

The Caesar cipher is defined by the following:

- key,  $k \in \mathbb{N}$ ,
- encryption function, e(x),
- decryption function, d(y),

$$e(x) = x + k \pmod{26},$$
  $d(y) = y - k \pmod{26}.$ 

The Caesar cipher is defined by the following:

- key,  $k \in \mathbb{N}$ ,
- encryption function, e(x),
- decryption function, d(y),

$$e(x) = x + k \pmod{26},$$
  $d(y) = y - k \pmod{26}.$ 

We first associate each letter to a number in  $\mathbb{Z}_{26}=\{0,1,\ldots,24,25\}$ , as follows:

A B C D E F G H I J K L M N O P Q R S T U V W X Y Z O 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25

The Caesar cipher is defined by the following:

- key,  $k \in \mathbb{N}$ ,
- $\blacksquare$  encryption function, e(x),
- decryption function, d(y),

$$e(x) = x + k \pmod{26},$$
  $d(y) = y - k \pmod{26}.$ 

We first associate each letter to a number in  $\mathbb{Z}_{26}=\{0,1,\ldots,24,25\}$ , as follows:

A B C D E F G H I J K L M N O P Q R S T U V W X Y Z O 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25

To see an example of this, suppose that k = 18.

The Caesar cipher is defined by the following:

- key,  $k \in \mathbb{N}$ ,
- encryption function, e(x),
- **decryption** function, d(y),

$$e(x) = x + k \pmod{26},$$
  $d(y) = y - k \pmod{26}.$ 

We first associate each letter to a number in  $\mathbb{Z}_{26}=\{0,1,\ldots,24,25\}$ , as follows:

To see an example of this, suppose that k = 18.

We encrypt the letter  ${f R}$  by

$$e(17) \equiv 17 + 18 \pmod{26}$$

The Caesar cipher is defined by the following:

- key,  $k \in \mathbb{N}$ ,
- $\blacksquare$  encryption function, e(x),
- **decryption** function, d(y),

$$e(x) = x + k \pmod{26},$$
  $d(y) = y - k \pmod{26}.$ 

We first associate each letter to a number in  $\mathbb{Z}_{26} = \{0,1,\dots,24,25\}$ , as follows:

To see an example of this, suppose that k = 18.

We encrypt the letter  ${f R}$  by

$$e(17) \equiv 17 + 18 \pmod{26}$$
  
 $\equiv 35 \pmod{26}$ 

The Caesar cipher is defined by the following:

- key,  $k \in \mathbb{N}$ ,
- $\blacksquare$  encryption function, e(x),
- **decryption** function, d(y),

$$e(x) = x + k \pmod{26},$$
  $d(y) = y - k \pmod{26}.$ 

We first associate each letter to a number in  $\mathbb{Z}_{26}=\{0,1,\ldots,24,25\}$ , as follows:

To see an example of this, suppose that k = 18.

We encrypt the letter R by

$$e(17) \equiv 17 + 18 \pmod{26}$$
  
 $\equiv 35 \pmod{26}$   
 $\equiv 9 \pmod{26}$ 

The Caesar cipher is defined by the following:

- key,  $k \in \mathbb{N}$ ,
- encryption function, e(x),
- **decryption** function, d(y),

$$e(x) = x + k \pmod{26},$$
  $d(y) = y - k \pmod{26}.$ 

We first associate each letter to a number in  $\mathbb{Z}_{26}=\{0,1,\ldots,24,25\}$ , as follows:

To see an example of this, suppose that k = 18.

We encrypt the letter  ${f R}$  by

$$e(17) \equiv 17 + 18 \pmod{26}$$
  
 $\equiv 35 \pmod{26}$   
 $\equiv 9 \pmod{26}$ 

which is .I

The Caesar cipher is defined by the following:

- key,  $k \in \mathbb{N}$ ,
- encryption function, e(x),
- **decryption** function, d(y),

$$e(x) = x + k \pmod{26},$$
  $d(y) = y - k \pmod{26}.$ 

We first associate each letter to a number in  $\mathbb{Z}_{26}=\{0,1,\ldots,24,25\}$ , as follows:

To see an example of this, suppose that k = 18.

We encrypt the letter R by

Let's decrypt L:

$$e(17) \equiv 17 + 18 \pmod{26}$$
  
 $\equiv 35 \pmod{26}$   
 $\equiv 9 \pmod{26}$ 

which is J.

The Caesar cipher is defined by the following:

- key,  $k \in \mathbb{N}$ ,
- encryption function, e(x),
- decryption function, d(y),

$$e(x) = x + k \pmod{26},$$
  $d(y) = y - k \pmod{26}.$ 

We first associate each letter to a number in  $\mathbb{Z}_{26}=\{0,1,\ldots,24,25\}$ , as follows:

To see an example of this, suppose that k = 18.

We encrypt the letter R by

Let's decrypt L:

$$e(17) \equiv 17 + 18 \pmod{26}$$
  
 $\equiv 35 \pmod{26}$   
 $\equiv 9 \pmod{26}$ 

which is J

 $d(11) \equiv 11 - 18 \pmod{26}$ 

The Caesar cipher is defined by the following:

- key,  $k \in \mathbb{N}$ ,
- encryption function, e(x),
- decryption function, d(y),

$$e(x) = x + k \pmod{26},$$
  $d(y) = y - k \pmod{26}.$ 

We first associate each letter to a number in  $\mathbb{Z}_{26}=\{0,1,\ldots,24,25\}$ , as follows:

To see an example of this, suppose that k = 18.

We encrypt the letter R by

$$e(17) \equiv 17 + 18 \pmod{26}$$
  $d(11) \equiv 11 - 18 \pmod{26}$   
 $\equiv 35 \pmod{26}$   $\equiv -7 \pmod{26}$   
 $\equiv 9 \pmod{26}$ 

which is .I

### Caesar cipher

The Caesar cipher is defined by the following:

- key,  $k \in \mathbb{N}$ ,
- encryption function, e(x),
- decryption function, d(y),

$$e(x) = x + k \pmod{26},$$
  $d(y) = y - k \pmod{26}.$ 

We first associate each letter to a number in  $\mathbb{Z}_{26}=\{0,1,\ldots,24,25\}$ , as follows:

To see an example of this, suppose that k = 18.

We encrypt the letter R by

$$e(17) \equiv 17 + 18 \pmod{26}$$
  $\qquad d(11) \equiv 11 - 18 \pmod{26}$   
 $\equiv 35 \pmod{26}$   $\qquad \equiv -7 \pmod{26}$   
 $\equiv 9 \pmod{26}$   $\qquad \equiv 19 \pmod{26}$ 

which is J.

### Caesar cipher

The Caesar cipher is defined by the following:

- key,  $k \in \mathbb{N}$ ,
- encryption function, e(x),
- **decryption** function, d(y),

$$e(x) = x + k \pmod{26},$$
  $d(y) = y - k \pmod{26}.$ 

We first associate each letter to a number in  $\mathbb{Z}_{26} = \{0, 1, \dots, 24, 25\}$ , as follows:

To see an example of this, suppose that k = 18.

We encrypt the letter R by

$$e(17) \equiv 17 + 18 \pmod{26}$$
  $d(11) \equiv 11 - 18 \pmod{26}$   
 $\equiv 35 \pmod{26}$   $\equiv -7 \pmod{26}$   
 $\equiv 9 \pmod{26}$   $\equiv 19 \pmod{26}$ 

which is .I

which is T

Consider the following encryption function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x \pmod{26}.$$

Consider the following encryption function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x \pmod{26}.$$

This works because the function e is injective, and this is because gcd(26,5) = 1.

Consider the following encryption function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x \pmod{26}.$$

This works because the function e is injective, and this is because gcd(26,5) = 1.

Consider the following encryption function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x \pmod{26}.$$

This works because the function e is injective, and this is because gcd(26,5) = 1.

$$e(17) \equiv 5 \cdot 17 \pmod{26}$$

Consider the following encryption function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x \pmod{26}.$$

This works because the function e is injective, and this is because gcd(26,5) = 1.

$$e(17) \equiv 5 \cdot 17 \pmod{26}$$
$$\equiv 85 \pmod{26}$$

Consider the following encryption function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x \pmod{26}.$$

This works because the function e is injective, and this is because gcd(26,5) = 1.

$$e(17) \equiv 5 \cdot 17 \pmod{26}$$
  
 $\equiv 85 \pmod{26}$   
 $\equiv 7 \pmod{26}$ .

Consider the following encryption function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x \pmod{26}.$$

This works because the function e is injective, and this is because gcd(26,5) = 1.

Let's encrypt the letter R, which is x = 17:

$$e(17) \equiv 5 \cdot 17 \pmod{26}$$
  
 $\equiv 85 \pmod{26}$   
 $\equiv 7 \pmod{26}$ .

The decryption function is

$$d(x) = 21x \pmod{26}.$$

Consider the following encryption function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x \pmod{26}.$$

This works because the function e is injective, and this is because gcd(26,5) = 1.

Let's encrypt the letter R, which is x = 17:

$$e(17) \equiv 5 \cdot 17 \pmod{26}$$
  
 $\equiv 85 \pmod{26}$   
 $\equiv 7 \pmod{26}$ .

The decryption function is

$$d(x) = 21x \pmod{26}.$$

This works because in  $\mathbb{Z}_{26}$ , the multiplicative inverse of k = 5 is  $5^{-1} := 21$ :

Consider the following encryption function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x \pmod{26}.$$

This works because the function e is injective, and this is because gcd(26,5) = 1.

Let's encrypt the letter R, which is x = 17:

$$e(17) \equiv 5 \cdot 17 \pmod{26}$$
  
 $\equiv 85 \pmod{26}$   
 $\equiv 7 \pmod{26}$ .

The decryption function is

$$d(x) = 21x \pmod{26}.$$

This works because in  $\mathbb{Z}_{26}$ , the multiplicative inverse of k=5 is  $5^{-1}:=21$ :

$$5 \cdot 21 \equiv 105 \pmod{26}$$

Consider the following encryption function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x \pmod{26}.$$

This works because the function e is injective, and this is because gcd(26,5) = 1.

Let's encrypt the letter R, which is x = 17:

$$e(17) \equiv 5 \cdot 17 \pmod{26}$$
  
 $\equiv 85 \pmod{26}$   
 $\equiv 7 \pmod{26}$ .

The decryption function is

$$d(x) = 21x \pmod{26}.$$

This works because in  $\mathbb{Z}_{26}$ , the multiplicative inverse of k=5 is  $5^{-1}:=21$ :

$$5 \cdot 21 \equiv 105 \pmod{26}$$
  
 $\equiv 1 \pmod{26}$ .

Consider the following encryption function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x \pmod{26}.$$

This works because the function e is injective, and this is because gcd(26,5) = 1.

Let's encrypt the letter R, which is x = 17:

$$e(17) \equiv 5 \cdot 17 \pmod{26}$$
  
 $\equiv 85 \pmod{26}$   
 $\equiv 7 \pmod{26}$ .

The decryption function is

$$d(x) = 21x \pmod{26}.$$

This works because in  $\mathbb{Z}_{26}$ , the multiplicative inverse of k=5 is  $5^{-1}:=21$ :

$$5 \cdot 21 \equiv 105 \pmod{26}$$
$$\equiv 1 \pmod{26}.$$

## Number theory fact

A number  $k \in \mathbb{Z}_n$  has a multiplicative inverse iff gcd(n, k) = 1.

Consider the following encrypting function:

$$e \colon \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad \qquad e(x) = 5x + 3 \pmod{26}.$$

Consider the following encrypting function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x + 3 \pmod{26}.$$

In other words, given the input x, we:

1. multiply by 5

Consider the following encrypting function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x + 3 \pmod{26}.$$

In other words, given the input x, we:

- 1. multiply by 5
- 2. add 3.

Consider the following encrypting function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x + 3 \pmod{26}.$$

In other words, given the input x, we:

- 1. multiply by 5
- 2. add 3.

To decrypt a message m, we need to "undo these" in the opposite order:

Consider the following encrypting function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x + 3 \pmod{26}.$$

In other words, given the input x, we:

- 1. multiply by 5
- 2. add 3.

To decrypt a message m, we need to "undo these" in the opposite order:

2. subtract 3

Consider the following encrypting function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x + 3 \pmod{26}.$$

In other words, given the input x, we:

- 1. multiply by 5
- 2. add 3.

To decrypt a message m, we need to "undo these" in the opposite order:

- 2. subtract 3
- 1. multiply by  $5^{-1} = 21$ .

Consider the following encrypting function:

$$e: \mathbb{Z}_{26} \to \mathbb{Z}_{26}, \qquad e(x) = 5x + 3 \pmod{26}.$$

In other words, given the input x, we:

- 1. multiply by 5
- 2. add 3.

To decrypt a message m, we need to "undo these" in the opposite order:

- 2. subtract 3
- 1. multiply by  $5^{-1} = 21$ .

The decryption function is thus

$$d(x) = 21(x-3) \pmod{26}$$
.

## A weakness of character ciphers

The ciphers that we've seen are called character, or monographic ciphers: all copies of the same letter get encrypted the same way:

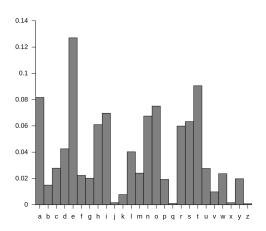
$$e(x_i) = e(x_j) \Rightarrow x_i = x_j.$$

# A weakness of character ciphers

The ciphers that we've seen are called character, or monographic ciphers: all copies of the same letter get encrypted the same way:

$$e(x_i) = e(x_j) \Rightarrow x_i = x_j.$$

If the message is long, then the private key can be deduced by analyzing letter frequencies.



A more sophisticated cipher are the block, or polygraphic ciphers, which encrypt blocks of plaintext letters to blocks of ciphertext letters of the same length.

A more sophisticated cipher are the block, or polygraphic ciphers, which encrypt blocks of plaintext letters to blocks of ciphertext letters of the same length.

One such system was developed by Blaise de Vigenère in 1585, called the Vigenère cipher.

A more sophisticated cipher are the block, or polygraphic ciphers, which encrypt blocks of plaintext letters to blocks of ciphertext letters of the same length.

One such system was developed by Blaise de Vigenère in 1585, called the Vigenère cipher.

We'll introduce this by an example.

Let's encrypt the message ENGINEERING:

$$p_1p_2p_3p_4p_5p_6p_7p_8p_9p_{10}p_{11} = 4 \ 13 \ 6 \ 8 \ 13 \ 4 \ 4 \ 17 \ 8 \ 13 \ 6$$

$$k_1k_2k_3k_4k_5 = 17 14 2 10 18.$$

A more sophisticated cipher are the block, or polygraphic ciphers, which encrypt blocks of plaintext letters to blocks of ciphertext letters of the same length.

One such system was developed by Blaise de Vigenère in 1585, called the Vigenère cipher.

We'll introduce this by an example.

Let's encrypt the message ENGINEERING:

$$p_1p_2p_3p_4p_5p_6p_7p_8p_9p_{10}p_{11} = 4 \ 13 \ 6 \ 8 \ 13 \ 4 \ 4 \ 17 \ 8 \ 13 \ 6$$

using the key ROCKS:

$$k_1 k_2 k_3 k_4 k_5 = 17 14 2 10 18.$$

ENGINEERING

A more sophisticated cipher are the block, or polygraphic ciphers, which encrypt blocks of plaintext letters to blocks of ciphertext letters of the same length.

One such system was developed by Blaise de Vigenère in 1585, called the Vigenère cipher.

We'll introduce this by an example.

Let's encrypt the message ENGINEERING:

$$p_1p_2p_3p_4p_5p_6p_7p_8p_9p_{10}p_{11} = 4 \ 13 \ 6 \ 8 \ 13 \ 4 \ 4 \ 17 \ 8 \ 13 \ 6$$

$$k_1 k_2 k_3 k_4 k_5 = 17 14 2 10 18.$$

A more sophisticated cipher are the block, or polygraphic ciphers, which encrypt blocks of plaintext letters to blocks of ciphertext letters of the same length.

One such system was developed by Blaise de Vigenère in 1585, called the Vigenère cipher.

We'll introduce this by an example.

Let's encrypt the message ENGINEERING:

$$p_1p_2p_3p_4p_5p_6p_7p_8p_9p_{10}p_{11} = 4 \ 13 \ 6 \ 8 \ 13 \ 4 \ 4 \ 17 \ 8 \ 13 \ 6$$

$$k_1 k_2 k_3 k_4 k_5 = 17 14 2 10 18.$$

A more sophisticated cipher are the block, or polygraphic ciphers, which encrypt blocks of plaintext letters to blocks of ciphertext letters of the same length.

One such system was developed by Blaise de Vigenère in 1585, called the Vigenère cipher.

We'll introduce this by an example.

Let's encrypt the message ENGINEERING:

$$p_1p_2p_3p_4p_5p_6p_7p_8p_9p_{10}p_{11} = 4 \ 13 \ 6 \ 8 \ 13 \ 4 \ 4 \ 17 \ 8 \ 13 \ 6$$

$$k_1 k_2 k_3 k_4 k_5 = 17 14 2 10 18.$$

A more sophisticated cipher are the block, or polygraphic ciphers, which encrypt blocks of plaintext letters to blocks of ciphertext letters of the same length.

One such system was developed by Blaise de Vigenère in 1585, called the Vigenère cipher.

We'll introduce this by an example.

Let's encrypt the message ENGINEERING:

$$p_1p_2p_3p_4p_5p_6p_7p_8p_9p_{10}p_{11} = 4 \ 13 \ 6 \ 8 \ 13 \ 4 \ 4 \ 17 \ 8 \ 13 \ 6$$

$$k_1k_2k_3k_4k_5 = 17 14 2 10 18.$$

A B C D E F G H I J K L M N O P Q R S T U V W X Y Z
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25

Let's decrypt the message TZGWK FBVSY WFU:

 $c_1c_2c_3c_4c_5c_6c_7c_8c_9c_{10}c_{11}c_{12}c_{13} = 19$  25 6 22 10 5 1 21 18 24 22 5 20

using the same key ROCKS:

 $k_1k_2k_3k_4k_5 = 17 14 2 10 18.$ 

A B C D E F G H I J K L M N O P Q R S T U V W X Y Z
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25

Let's decrypt the message TZGWK FBVSY WFU:

 $c_1c_2c_3c_4c_5c_6c_7c_8c_9c_{10}c_{11}c_{12}c_{13} = 19$  25 6 22 10 5 1 21 18 24 22 5 20

using the same key  $\mathsf{ROCKS}$ :

 $k_1k_2k_3k_4k_5 = 17 14 2 10 18.$ 

T Z G W K F B V S Y W F U

Let's decrypt the message TZGWK FBVSY WFU:

 $c_1c_2c_3c_4c_5c_6c_7c_8c_9c_{10}c_{11}c_{12}c_{13}=19\ 25\ 6\ 22\ 10\ 5\ 1\ 21\ 18\ 24\ 22\ 5\ 20$  using the same key **ROCKS**:

 $k_1k_2k_3k_4k_5 = 17 14 2 10 18.$ 

Let's decrypt the message TZGWK FBVSY WFU:

 $c_1c_2c_3c_4c_5c_6c_7c_8c_9c_{10}c_{11}c_{12}c_{13}=19\ \ 25\ \ 6\ \ 22\ \ 10\ \ 5\ \ 1\ \ 21\ \ 18\ \ 24\ \ 22\ \ 5\ \ 20$  using the same key **ROCKS**:

 $k_1k_2k_3k_4k_5 = 17 14 2 10 18.$ 

	Т												
$c_i$	19	25	6	22	10	5	1	21	18	24	22	5	20
$k_i$	17	14	2	10	18	17	14	2	10	18	17	14	2

#### Let's decrypt the message TZGWK FBVSY WFU:

 $c_1c_2c_3c_4c_5c_6c_7c_8c_9c_{10}c_{11}c_{12}c_{13}=19\ 25\ 6\ 22\ 10\ 5\ 1\ 21\ 18\ 24\ 22\ 5\ 20$  using the same key **ROCKS**:

$$k_1k_2k_3k_4k_5 = 17 14 2 10 18.$$

	Т	Z	G	W	K	F	В	V	S	Υ	W	F	U
$c_i$				22									
				10									
$p_i = c_i - k_i$	2	11	4	12	18	14	13	19	8	6	4	17	18

#### Let's decrypt the message TZGWK FBVSY WFU:

 $c_1c_2c_3c_4c_5c_6c_7c_8c_9c_{10}c_{11}c_{12}c_{13}=19\ \ 25\ \ 6\ \ 22\ \ 10\ \ 5\ \ 1\ \ 21\ \ 18\ \ 24\ \ 22\ \ 5\ \ 20$  using the same key **ROCKS**:

 $k_1k_2k_3k_4k_5 = 17 14 2 10 18.$ 

	Т												
	19												
$k_i$													
$p_i = c_i - k_i$													
	С	L	Ε	М	S	0	N	Т	ı	G	Ε	R	S

The ciphers in this lecture are symmetric: Decryption is the opposite ("inverse") of encryption.

The ciphers in this lecture are symmetric: Decryption is the opposite ("inverse") of encryption.

Not only is the same key is used for encryption and decryption, but that key needs to be kept private.

The ciphers in this lecture are symmetric: Decryption is the opposite ("inverse") of encryption.

Not only is the same key is used for encryption and decryption, but that key needs to be kept private.

This is a serious problem if two entities want to send a message across an insecure channel, e.g., a customer and an online bank.

The ciphers in this lecture are symmetric: Decryption is the opposite ("inverse") of encryption.

Not only is the same key is used for encryption and decryption, but that key needs to be kept private.

This is a serious problem if two entities want to send a message across an insecure channel, e.g., a customer and an online bank.

However, there are secure symmetric cryptosystems like the popular Diffie–Hellman (DH) key exchange.

The ciphers in this lecture are symmetric: Decryption is the opposite ("inverse") of encryption.

Not only is the same key is used for encryption and decryption, but that key needs to be kept private.

This is a serious problem if two entities want to send a message across an insecure channel, e.g., a customer and an online bank.

However, there are secure symmetric cryptosystems like the popular Diffie–Hellman (DH) key exchange.

#### Definition

In an asymmetric cipher, there are two distinct keys:

A public key, used for encryption;

The ciphers in this lecture are symmetric: Decryption is the opposite ("inverse") of encryption.

Not only is the same key is used for encryption and decryption, but that key needs to be kept private.

This is a serious problem if two entities want to send a message across an insecure channel, e.g., a customer and an online bank.

However, there are secure symmetric cryptosystems like the popular Diffie–Hellman (DH) key exchange.

#### Definition

In an asymmetric cipher, there are two distinct keys:

- A public key, used for encryption;
- A private key, used for decryption.

The ciphers in this lecture are symmetric: Decryption is the opposite ("inverse") of encryption.

Not only is the same key is used for encryption and decryption, but that key needs to be kept private.

This is a serious problem if two entities want to send a message across an insecure channel, e.g., a customer and an online bank.

However, there are secure symmetric cryptosystems like the popular Diffie–Hellman (DH) key exchange.

#### Definition

In an asymmetric cipher, there are two distinct keys:

- A public key, used for encryption;
- A private key, used for decryption.

The instructions for encrypting a message can be made public without compromising the security.

The ciphers in this lecture are symmetric: Decryption is the opposite ("inverse") of encryption.

Not only is the same key is used for encryption and decryption, but that key needs to be kept private.

This is a serious problem if two entities want to send a message across an insecure channel, e.g., a customer and an online bank.

However, there are secure symmetric cryptosystems like the popular Diffie–Hellman (DH) key exchange.

#### Definition

In an asymmetric cipher, there are two distinct keys:

- A public key, used for encryption;
- A private key, used for decryption.

The instructions for encrypting a message can be made public without compromising the security.

The Rivest–Shamir–Adleman (RSA) cryptosystem is a popular asymmetric cryptosystem, which we'll learn about this in an upcoming lecture.