

Lecture 4.3: Partially ordered sets

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Math 4190, Discrete Mathematical Structures

Recall the basic concepts

Definition

A **partial order** on a set P is a relation \preceq that is

- (i) reflexive,
- (ii) transitive,
- (iii) antisymmetric.

We say that (P, \preceq) is a **partially ordered set**, or a **poset**.

Definition (alternate)

A (strict) **partial order** on a set P is a relation \prec that is

- (i) irreflexive,
- (ii) transitive,
- (iii) antisymmetric

Definition

If (P, \preceq) is a poset, and for every $a \neq b \in P$, either $a \preceq b$, or $b \preceq a$, then it is a **totally ordered set**.

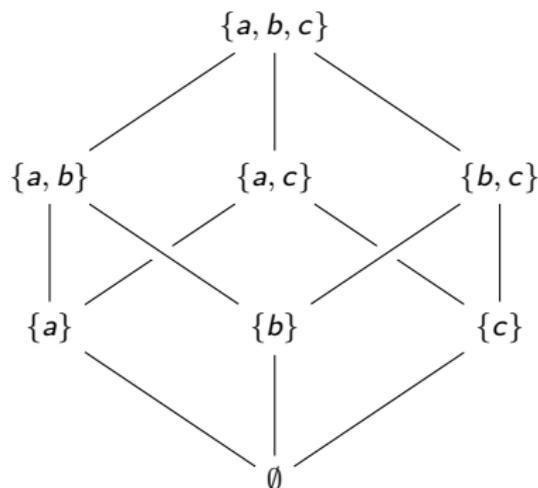
A bunch of definitions

Let (P, \preceq) be a poset, and $a, x, y, z \in P$.

- If $x \not\preceq y$ and $y \not\preceq x$, then x and y are **incomparable**. Otherwise, they are **comparable**.
- If $x \preceq z$, but $\nexists y \in P$ such that $x \prec y \prec z$, then z **covers** x .
- If $a \in P$ but $\nexists x \in P$ such that $x \prec a$, then a is a **minimal element**.
- If $a \preceq x$ for all $x \in P$, then a is the **minimum** element.
- If $z \in P$ but $\nexists x \in P$ such that $z \prec x$, then z is a **maximal element**.
- If $x \preceq z$ for all $x \in P$, then z is the **maximum** element.
- A **chain** in a poset is a subset $C \subseteq P$ such that any two elements are comparable.
- An **antichain** in a poset is a subset $A \subseteq P$ of incomparable elements.
- A poset $(P, \leq_{P'})$ is an **extension** of (P, \leq_P) if $P = P'$ and $a \leq_P b$ implies $a \leq_{P'} b$.
- A **linear extension** of a poset is an extension that is a total order.

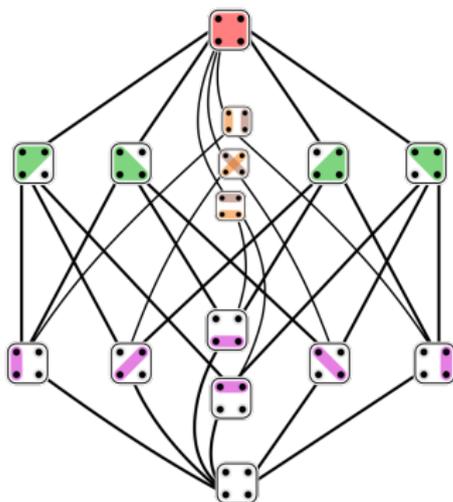
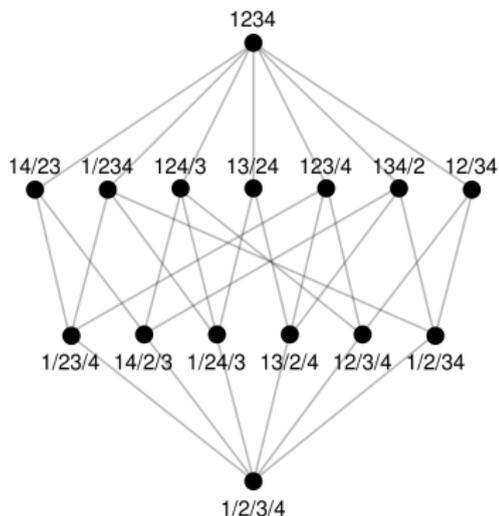
Examples of posets

1. Power set: $(2^S, \subseteq)$.
2. Partitions of $[n] = \{1, \dots, n\}$.
3. Any acyclic directed graph.
4. Divisors of n . Or the integers, by divisibility: $(\mathbb{Z}^+, |)$.
5. Vertices in a rooted tree (e.g., computer directory structure, phylogenetic tree).
6. Strongly connected components in a directed graph.



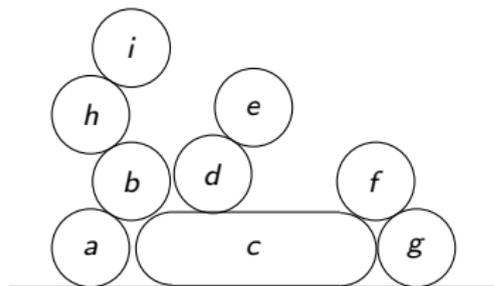
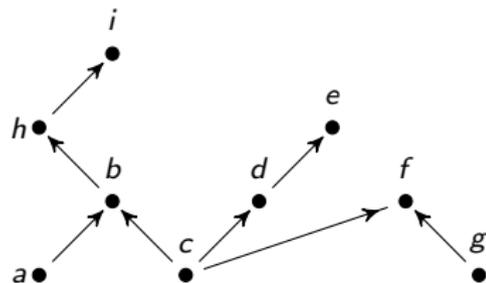
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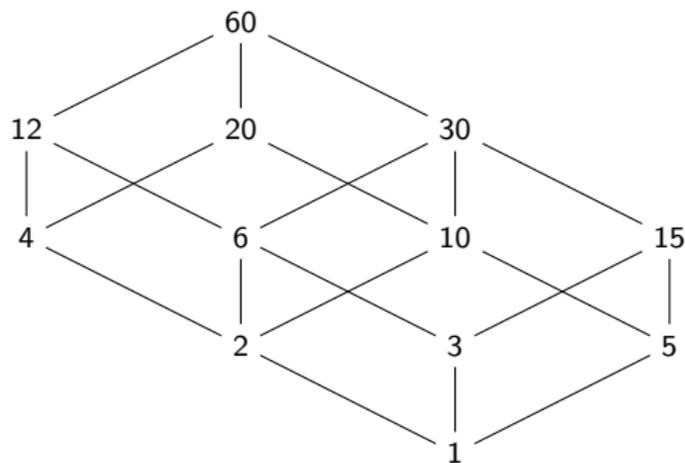
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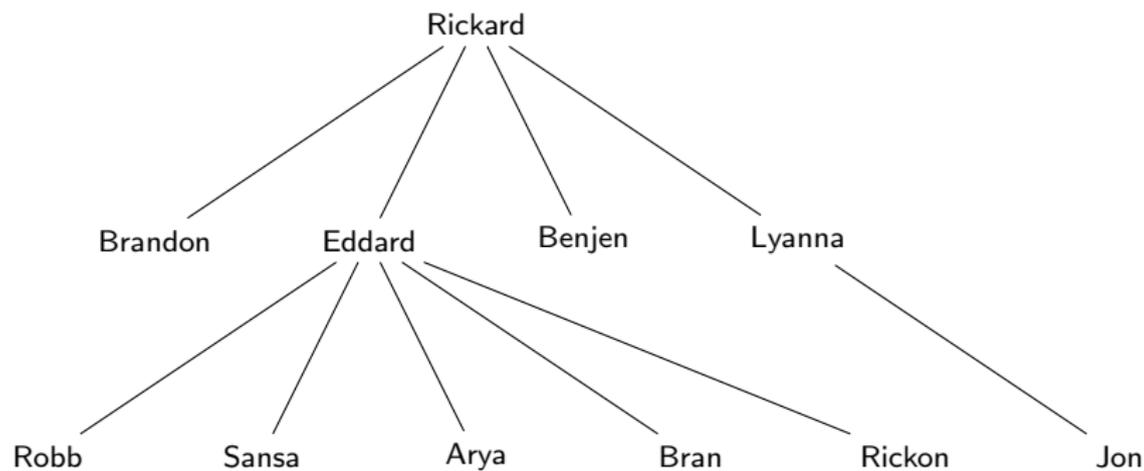
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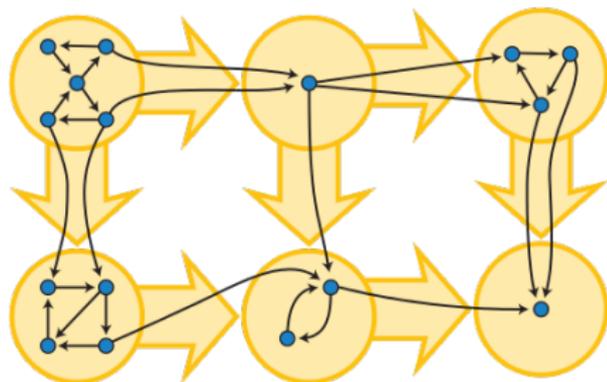
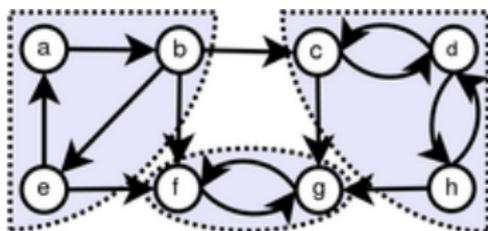
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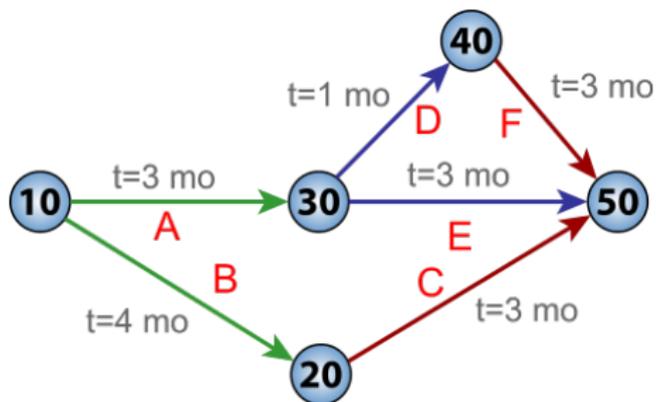


Applications to scheduling problems

The field of **operations research** deals with methods (algorithms, optimization, heuristics, etc.) that improve complex decision making.

In the 1950s, the **Program Evaluation and Review Technique (PERT)** was developed by the U.S. Navy when they were building the Polaris submarine.

Around this time, the **Critical Path Method (CPM)** was developed by the DuPont chemical company for scheduling maintenance. It was later used during the construction of the World Trade Center.



Applications to scheduling: PERT and CPM

In both PERT and CPM, the set of scheduled tasks forms a **partially ordered set**.

The tasks are labeled with the duration that they take to complete.

The shortest possible completion time is given by a maximal chain called a **critical path**.

