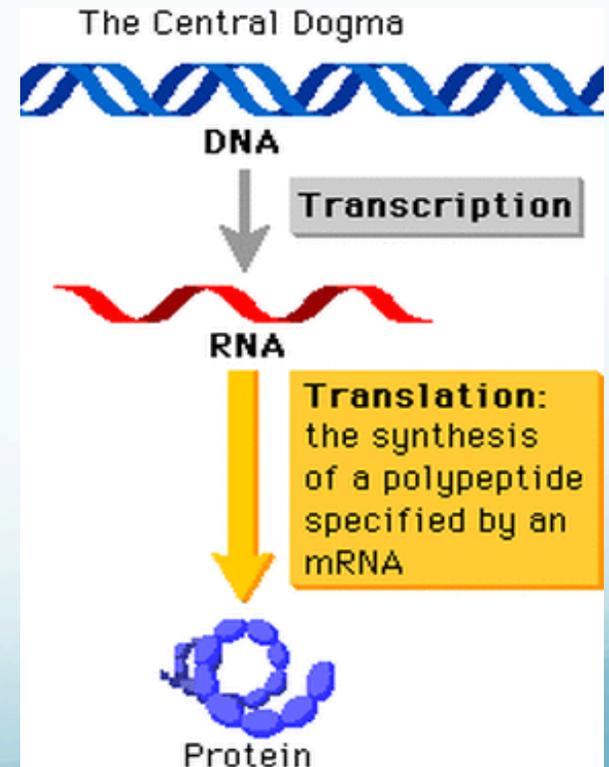


Boolean models of the *lac* operon in *E. coli*

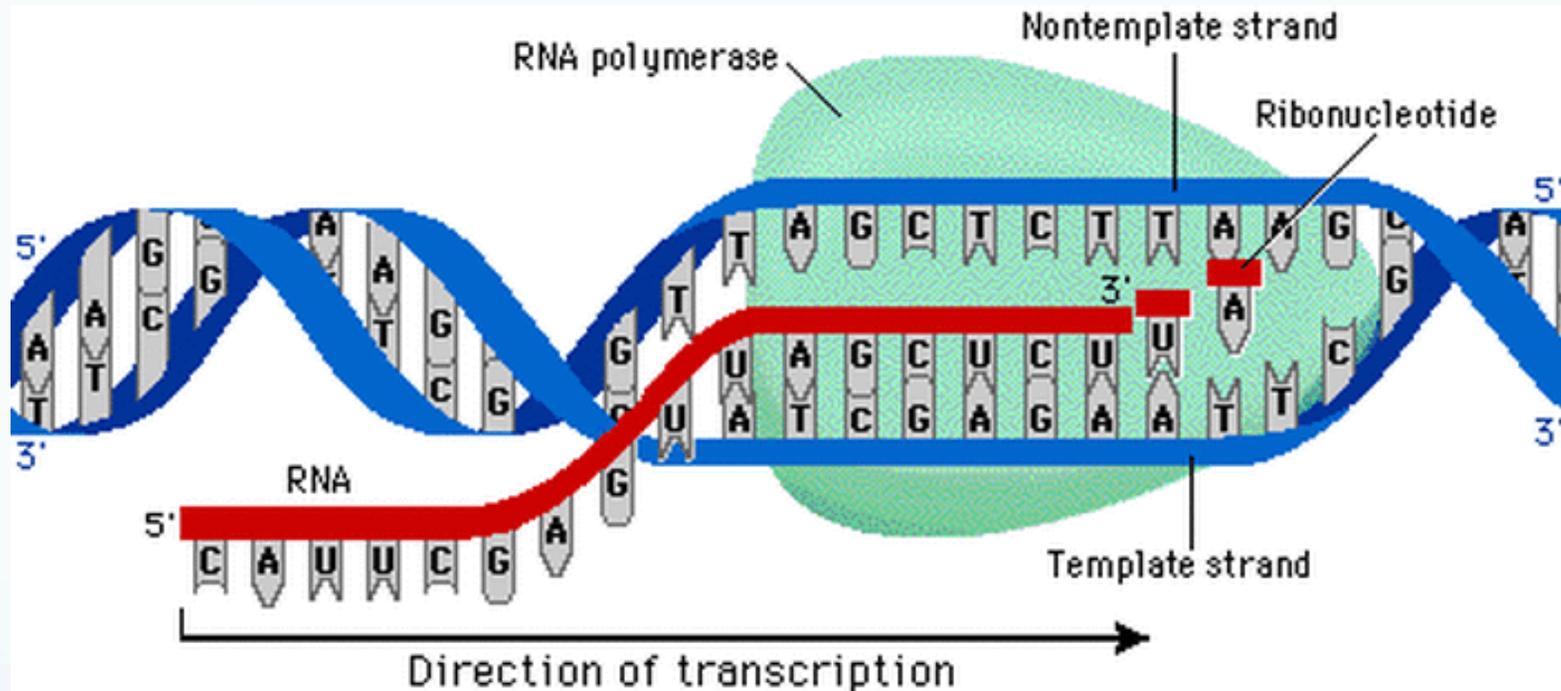
Matthew Macauley
Clemson University

Gene expression

- **Gene expression** is a process that takes gene info and creates a functional gene product (e.g., a protein).
- Some genes code for proteins. Others (e.g., rRNA, tRNA) code for functional RNA.
- Gene Expression is a 2-step process:
 - 1) **transcription** of genes (messenger RNA synthesis)
 - 2) **translation** of genes (protein synthesis)
- DNA consists of bases A, C, G, T.
- RNA consists of bases A, C, G, U.
- Proteins are long chains of amino acids.
- Gene expression is used by all known life forms.

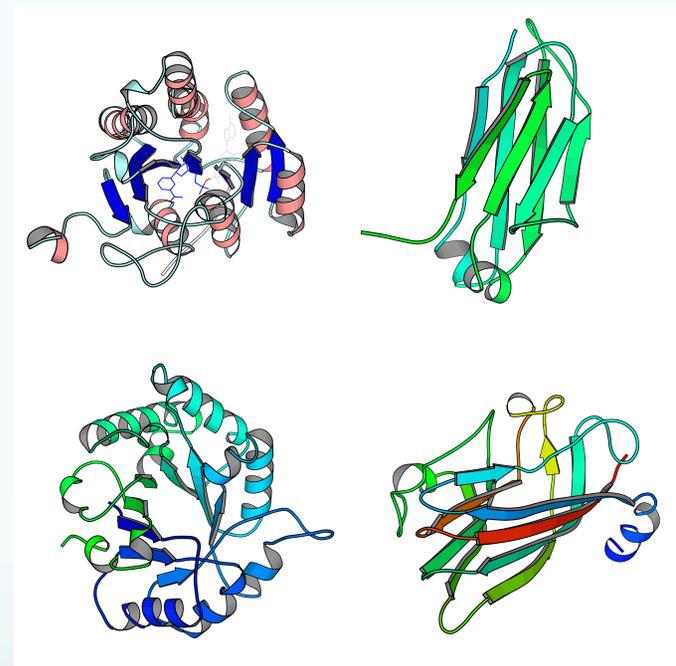
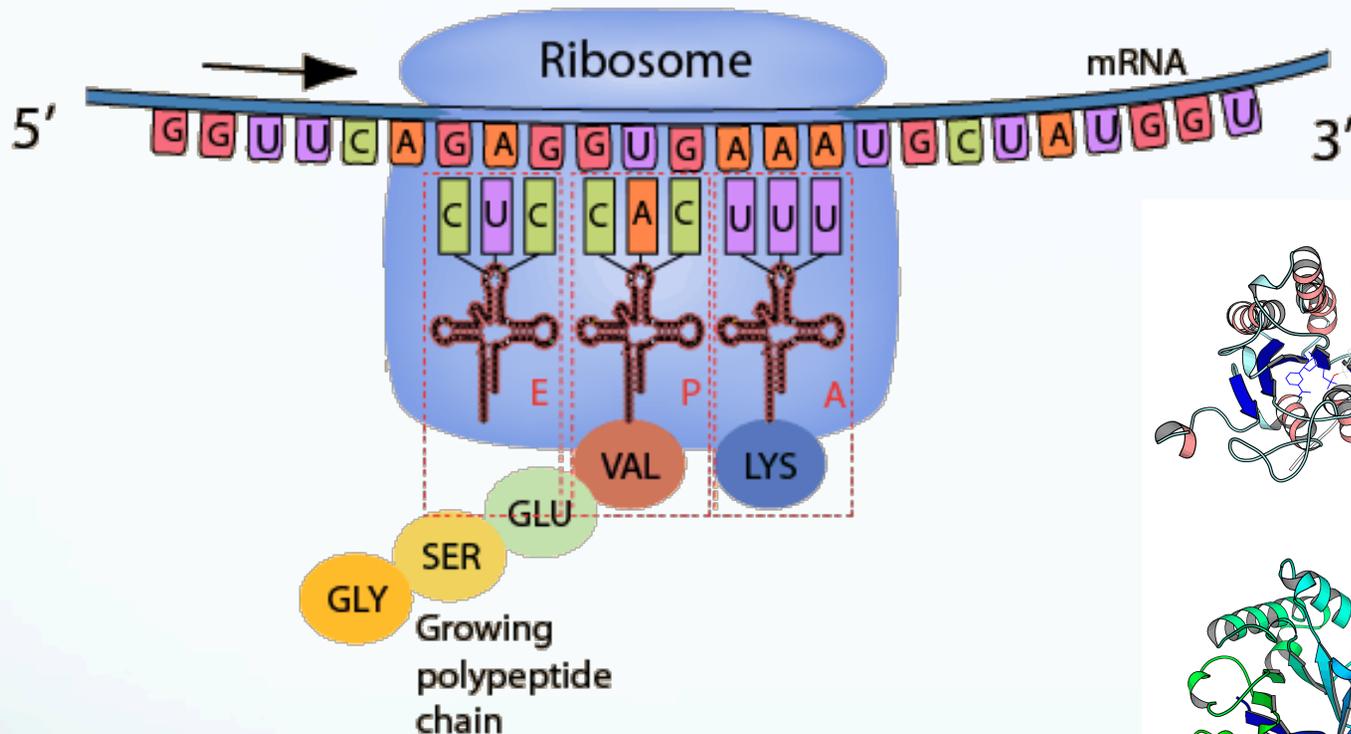


Transcription



- Transcription occurs inside the cell nucleus.
- A helicase enzyme binds to and “unzips” DNA to read it.
- DNA is copied into mRNA.
- Segments of RNA not needed for protein coding are removed.
- The RNA then leaves the cell nucleus.

Translation



- During translation, the mRNA is read by **ribosomes**.
- Each triple of RNA bases codes for an **amino acid**.
- The result is a **protein**: a long chain of amino acids.
- Proteins fold into a 3-D shape which determine their function

Gene expression

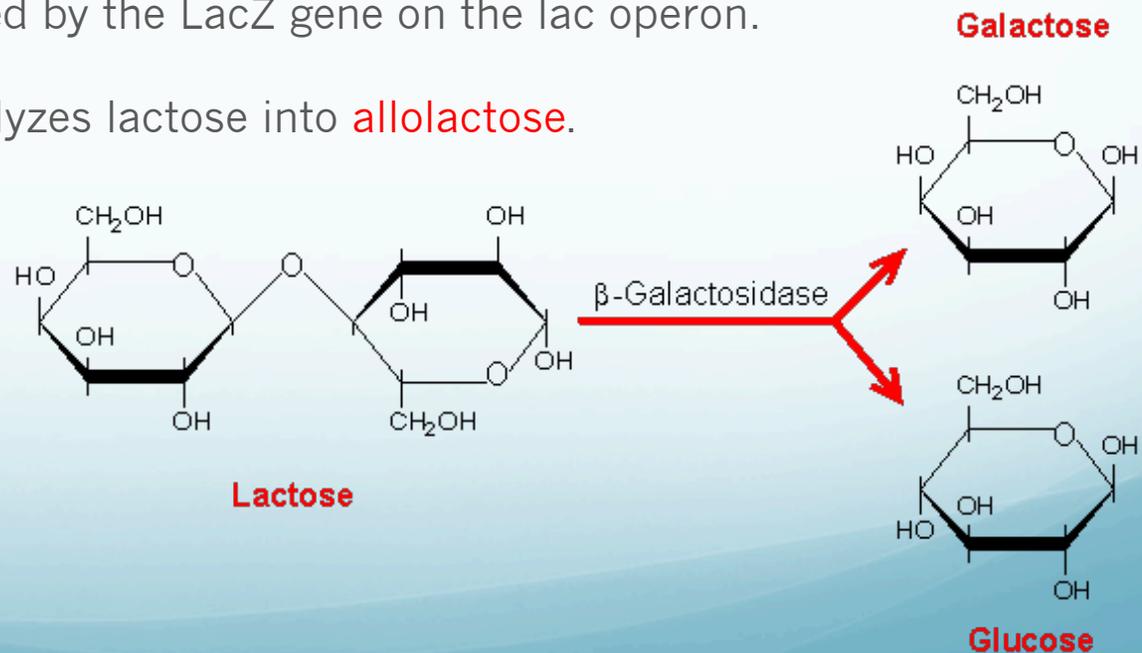
- The **expression level** is the **rate** at which a gene is being expressed.
- **Housekeeping genes** are continuously expressed, as they are essential for basic life processes.
- **Regulated genes** are expressed only under certain outside factors (environmental, physiological, etc.). Expression is controlled by the cell.
- It is easiest to control gene regulation by affecting transcription.
- One way to block repression is for **repressor proteins** bind to the DNA or RNA.
- **Goal**: Understand the complex cell behaviors of **gene regulation**, which is the process of turning on/off certain genes depending on the requirements of the organism.

The *lac* operon in *E. coli*

- An **operon** is a region of DNA that contains a cluster of genes that are transcribed together.
- *E. coli* is a bacterium in the gut of mammals and birds. Its genome has been sequenced and its physiology is well-understood.
- The **lactose (*lac*) operon** controls the **transport** and **metabolism** of lactose in *Escherichia coli*.
- The *lac* operon was discovered by Francois Jacob and Jacques Monod in 1961, which earned them the Nobel Prize.
- The *lac* operon was the first operon discovered and is the most widely studied mechanism of gene regulation.
- The *lac* operon is used as a “test system” for models of gene regulation.
- DNA replication and gene expression were all studied in *E. coli* before they were studied in eukaryotic cells.

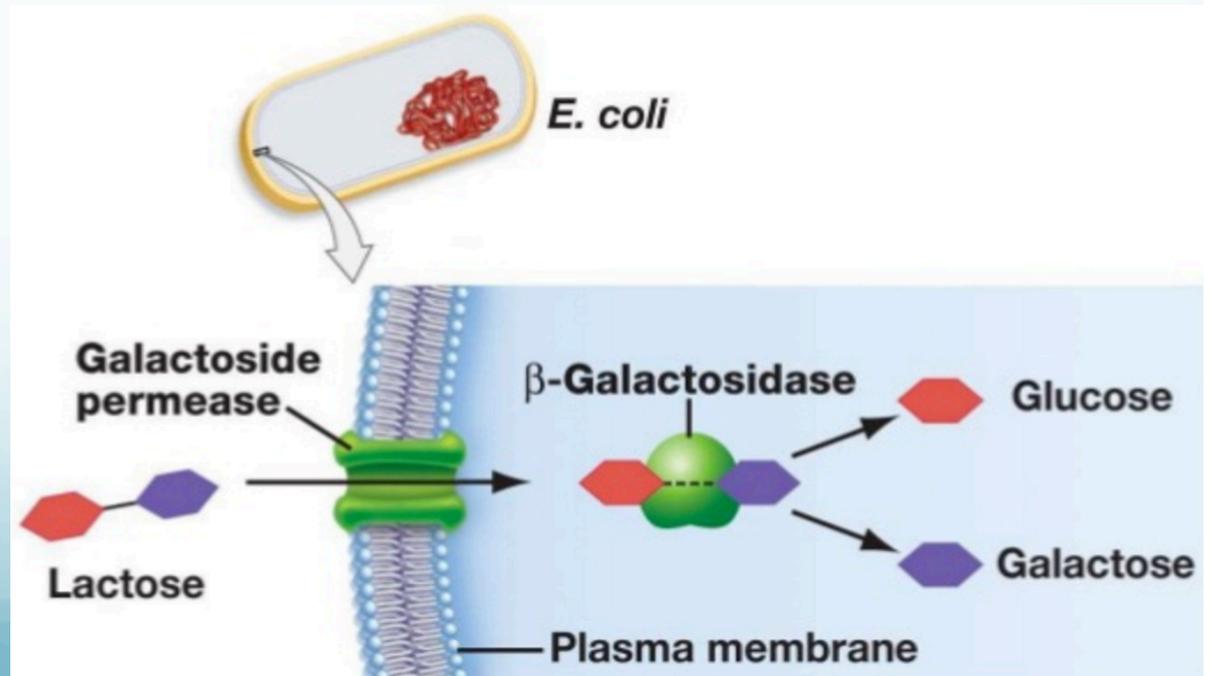
Lactose and β -galactosidase

- When a host consumes milk, *E. coli* is exposed to **lactose** (milk sugar).
- Lactose consists of one **glucose sugar** linked to one **galactose sugar**.
- If both glucose and lactose are available, then glucose is the preferred energy source.
- Before lactose can be used as energy, the **β -galactosidase** enzyme is needed to break it down.
- β -galactosidase is encoded by the LacZ gene on the lac operon.
- β -galactosidase also catalyzes lactose into **allolactose**.

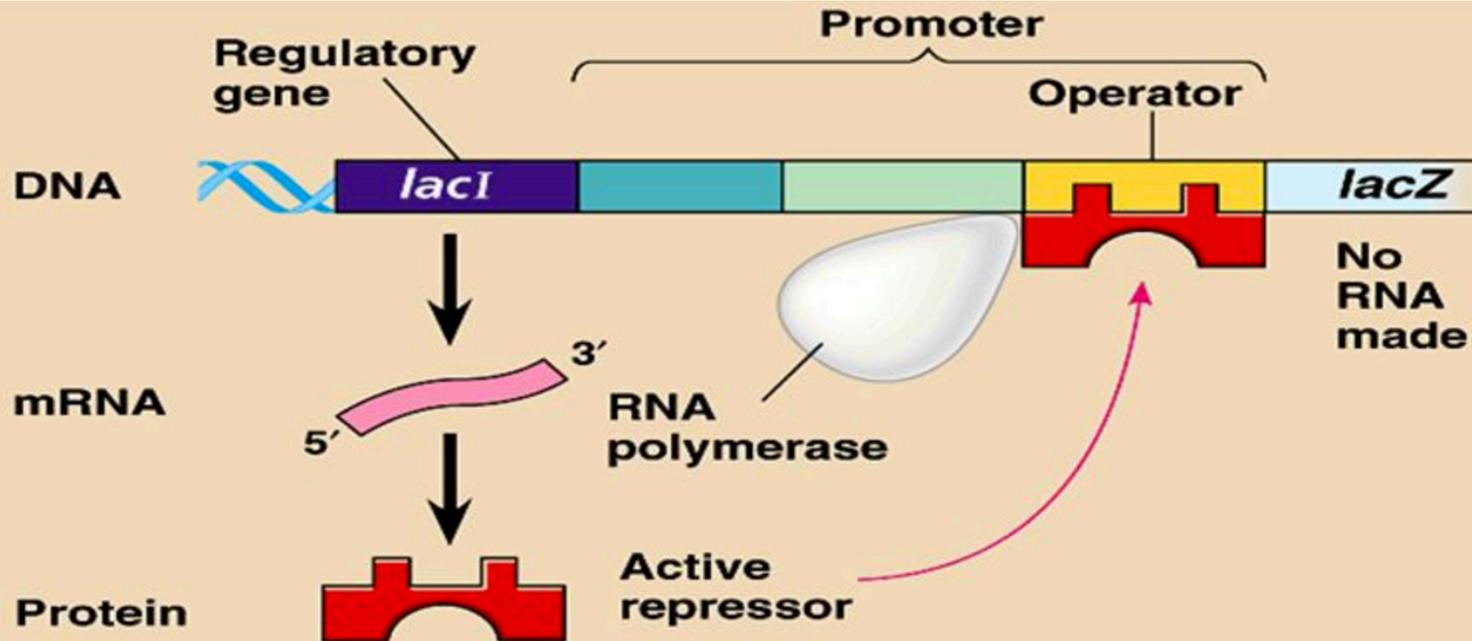


Transporter protein

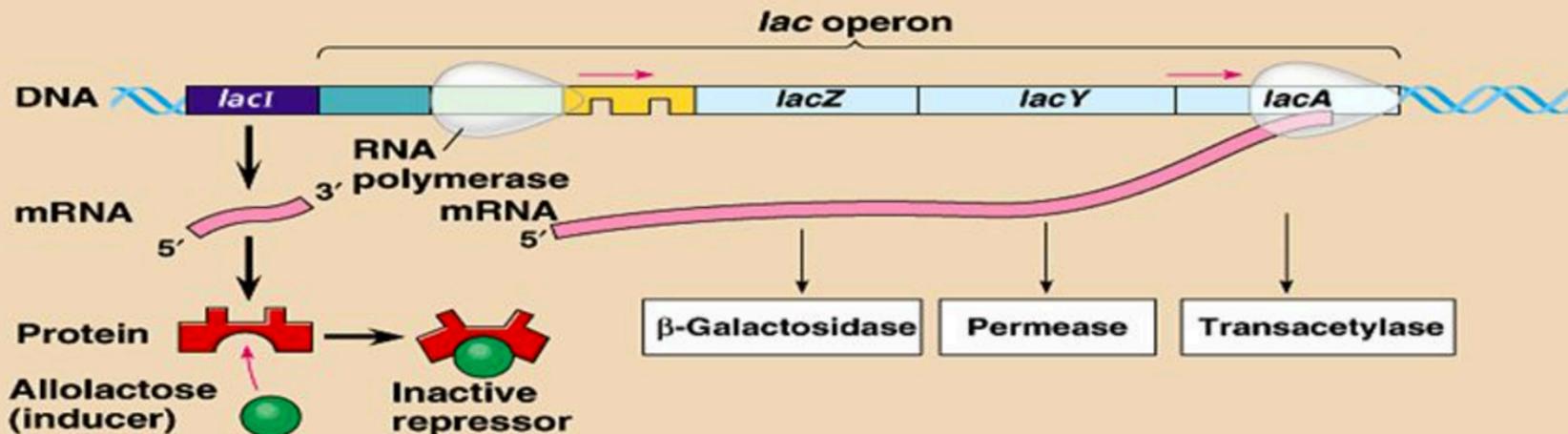
- To bring lactose into the cell, a **transport protein**, called *lac permease*, is required.
- This protein is encoded by the LacY gene on the *lac* operon.
- If lactose is not present, then neither of the following are produced:
 - 1) β -galactosidase (LacZ gene)
 - 2) *lac* permease (LacY gene)
- In this case, the *lac* operon is OFF.



The *lac* operon



(a) Lactose absent, repressor active, operon off



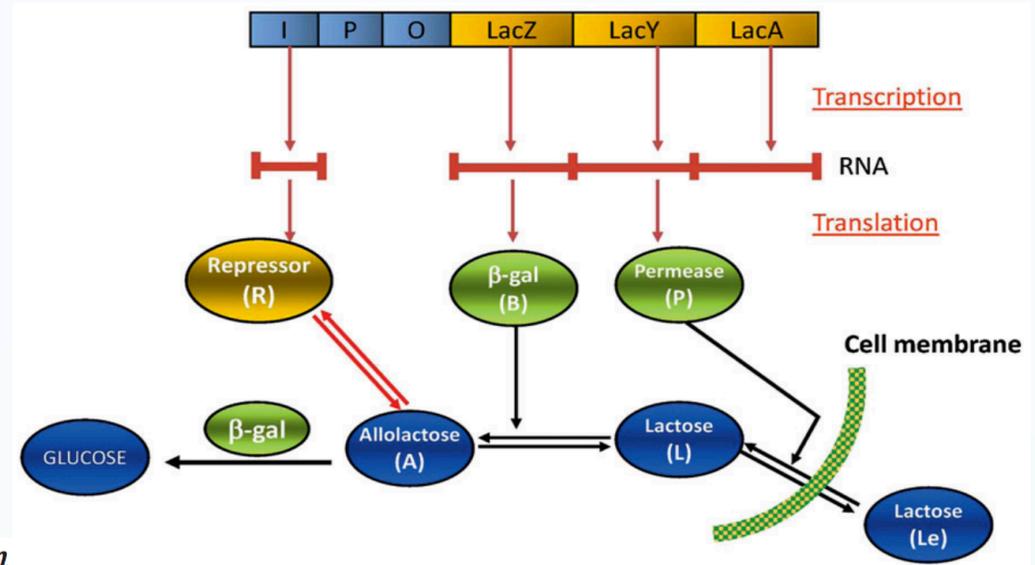
(b) Lactose present, repressor inactive, operon on

with lactose and no glucose

- Lactose is brought into the cell by the *lac* permease transporter protein
- β -galactosidase breaks up lactose into glucose and galactose..
- β -galactosidase also converts lactose into allolactose.
- Allolactose binds to the *lac* repressor protein, preventing it from binding to the operator region of the genome.
- Transcription begins: mRNA encoding the *lac* genes is produced.
- Lac proteins are produced, and more lactose is brought into the cell. (The operon is ON.)
- Eventually, all lactose is used up, so there will be no more allolactose.
- The *lac* repressor can now bind to the operator, so mRNA transcription stops. (The operon has turned itself OFF.)

An ODE *lac* operon model

- M: mRNA
- B: β -galactosidase
- A: allolactose
- P: transporter protein
- L: lactose



$$\frac{dM}{dt} = \alpha_M \frac{1 + K_1 (e^{-\mu\tau_M} A\tau_M)^n}{K + K_1 (e^{-\mu\tau_M} A\tau_M)^n} + \Gamma_0 - \tilde{\gamma}_M M$$

$$\frac{dB}{dt} = \alpha_B e^{-\mu\tau_B} M\tau_B - \tilde{\gamma}_B B$$

$$\frac{dA}{dt} = \alpha_A B \frac{L}{K_L + L} - \beta_A B \frac{A}{K_A + A} - \tilde{\gamma}_A A$$

$$\frac{dP}{dt} = \alpha_P e^{-\mu(\tau_B + \tau_P)} M\tau_{B+\tau_P} - \tilde{\gamma}_P P$$

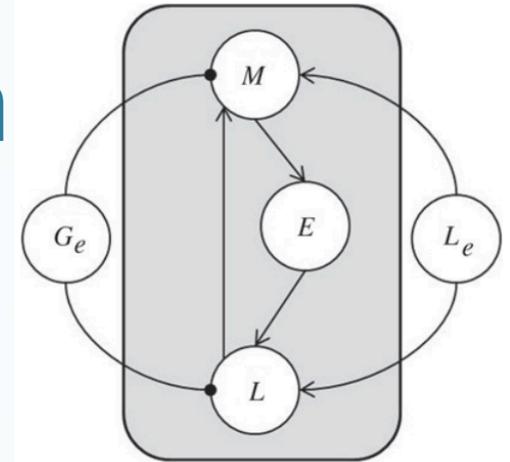
$$\frac{dL}{dt} = \alpha_L P \frac{L_e}{K_{L_e} + L_e} - \beta_{L_1} P \frac{L}{K_{L_1} + L} - \alpha_A B \frac{L}{K_L + L} - \tilde{\gamma}_L L$$

Downsides of an ODE model

- Very mathematically advanced.
- Too hard to solve explicitly. Numerical methods are needed.
- MANY experimentally determined “rate constants” (I count 18...)
- Often, these rate constants aren't known even up to orders of magnitude.

$$\begin{aligned}\frac{dM}{dt} &= \alpha_M \frac{1 + K_1 (e^{-\mu\tau_M} A_{\tau_M})^n}{K + K_1 (e^{-\mu\tau_M} A_{\tau_M})^n} + \Gamma_0 - \tilde{\gamma}_M M \\ \frac{dB}{dt} &= \alpha_B e^{-\mu\tau_B} M_{\tau_B} - \tilde{\gamma}_B B \\ \frac{dA}{dt} &= \alpha_A B \frac{L}{K_L + L} - \beta_A B \frac{A}{K_A + A} - \tilde{\gamma}_A A \\ \frac{dP}{dt} &= \alpha_P e^{-\mu(\tau_B + \tau_P)} M_{\tau_B + \tau_P} - \tilde{\gamma}_P P \\ \frac{dL}{dt} &= \alpha_L P \frac{L_e}{K_{L_e} + L_e} - \beta_{L_1} P \frac{L}{K_{L_1} + L} - \alpha_A B \frac{L}{K_L + L} - \tilde{\gamma}_L L\end{aligned}$$

A Boolean approach



- Let's assume everything is "Boolean" (0 or 1):
 - Gene products are either present or absent
 - Enzyme concentrations are either high or low.
 - The operon is either ON or OFF.
- mRNA is transcribed ($M=1$) if there is no external glucose ($G=0$), and either internal lactose ($L=1$) or external lactose ($L_e=1$) are present.

$$x_M(t+1) = f_M(t+1) = \overline{G_e} \wedge (L(t) \vee L_e)$$

- The LacY and LacZ gene products ($E=1$) will be produced if mRNA is available ($M=1$).

$$x_E(t+1) = f_E(t+1) = M(t)$$

- Lactose will be present in the cell if there is no external glucose ($G_e=0$), and either of the following holds:
 - ✓ External lactose is present ($L_e=1$) and *lac* permease ($E=1$) is available.
 - ✓ Internal lactose is present ($L=1$), but β -galactosidase is absent ($E=0$).

$$x_L(t+1) = f_L(t+1) = \overline{G_e} \wedge \left[(L_e \wedge E(t)) \vee (L(t) \wedge \overline{E(t)}) \right]$$

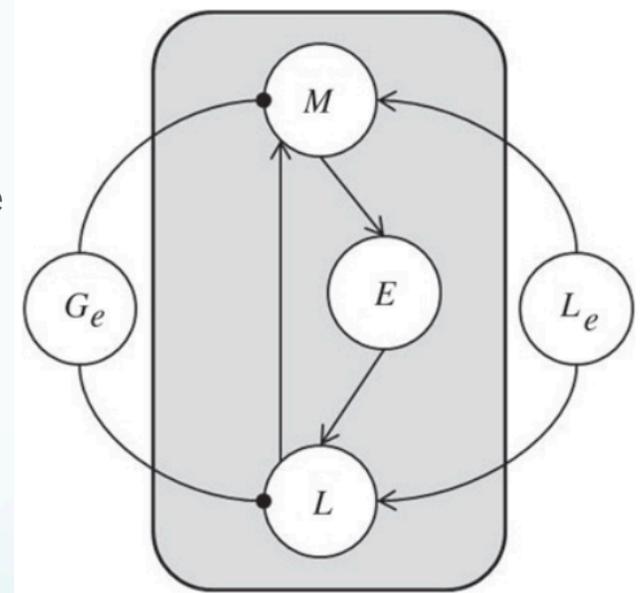
Comments on the Boolean model

- We have two “types” of Boolean quantities:
 - mRNA (M), lac gene products (E), and internal lactose (L) are **variables**.
 - External glucose (G_e) and lactose (L_e) are **parameters** (constants).
- Variables and parameters are drawn as **nodes**.
- Interactions can be drawn as **signed edges**.
- A signed graph called the **wiring diagram** describes the dependencies of the variables.
- Time is discrete: $t = 0, 1, 2, \dots$

$$x_M(t+1) = f_M(t+1) = \overline{G_e} \wedge (L(t) \vee L_e)$$

$$x_E(t+1) = f_E(t+1) = M(t)$$

$$x_L(t+1) = f_L(t+1) = \overline{G_e} \wedge \left[(L_e \wedge E(t)) \vee (L(t) \wedge \overline{E(t)}) \right]$$



- Assume that the variables are updated **synchronously**.

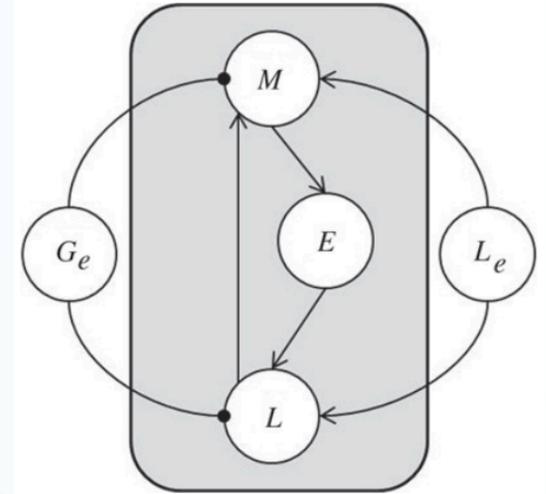
How to analyze a Boolean model

- At the bare minimum, we should expect:
 - Lactose absent => operon OFF.
 - Lactose present, glucose absent => operon ON.
 - Lactose and glucose present => operon OFF.

$$x_M(t+1) = f_M(t+1) = \overline{G_e} \wedge (L(t) \vee L_e)$$

$$x_E(t+1) = f_E(t+1) = M(t)$$

$$x_L(t+1) = f_L(t+1) = \overline{G_e} \wedge [(L_e \wedge E(t)) \vee (L(t) \wedge \overline{E(t)})]$$



- The state space (or phase space) is the directed graph (V, T) , where

$$V = \{(x_M, x_E, x_L) : x_i \in \{0,1\}\} \quad T = \{(x, f(x)) : x \in V\}$$

- We'll draw the state space for all four choices of the parameters:
 - $(L_e, G_e) = (0, 0)$. We hope to end up in a fixed point $(0,0,0)$.
 - $(L_e, G_e) = (0, 1)$. We hope to end up in a fixed point $(0,0,0)$.
 - $(L_e, G_e) = (1, 0)$. We hope to end up in a fixed point $(1,1,1)$.
 - $(L_e, G_e) = (1, 1)$. We hope to end up in a fixed point $(0,0,0)$.

How to analyze a Boolean model

- We can plot the state space using the software: Analysis of Dynamical Algebraic Models (ADAM), at adam.plantsimlab.org.
- First, we need to convert our logical functions into polynomials.

$$x_M(t+1) = f_M(t+1) = \overline{G_e} \wedge (L(t) \vee L_e)$$

$$x_E(t+1) = f_E(t+1) = M(t)$$

$$x_L(t+1) = f_L(t+1) = \overline{G_e} \wedge \left[(L_e \wedge E(t)) \vee (L(t) \wedge \overline{E(t)}) \right]$$

- Here is the relationship between Boolean logic and polynomial algebra:

| <u>Boolean operations</u> | <u>logical form</u> | <u>polynomial form</u> |
|---------------------------|---------------------|------------------------|
| ○ AND | $z = x \wedge y$ | $z = xy$ |
| ○ OR | $z = x \vee y$ | $z = x + y + xy$ |
| ○ NOT | $z = \overline{x}$ | $z = 1 + x$ |

- Also, everything is done modulo 2, so $1+1=0$, and $x^2=x$, and thus $x(x+1)=0$.

Analysis of Dynamic Algebraic Models (ADAM) v1.1

$$x_M(t+1) = f_M(t+1) = \overline{G_e} \wedge (L(t) \vee L_e)$$

$$x_E(t+1) = f_E(t+1) = M(t)$$

$$x_L(t+1) = f_L(t+1) = \overline{G_e} \wedge \left[(L_e \wedge E(t)) \vee (L(t) \wedge \overline{E(t)}) \right]$$

Model Input:

Upload a file (.txt) for the polynomials corresponding to the system OR enter them directly into the text area:

SELECT FILES

Enter external parameters directly into the text area; one for each line:

```
f1 = (1+Ge)*(x3*Le+x3+Le)
f2 = x1
f3 = (1+Ge)*(x2*Le+x3*(1+x2))
```

```
Ge = 0
Le = 1
```

4)

What analysis would you like to run?

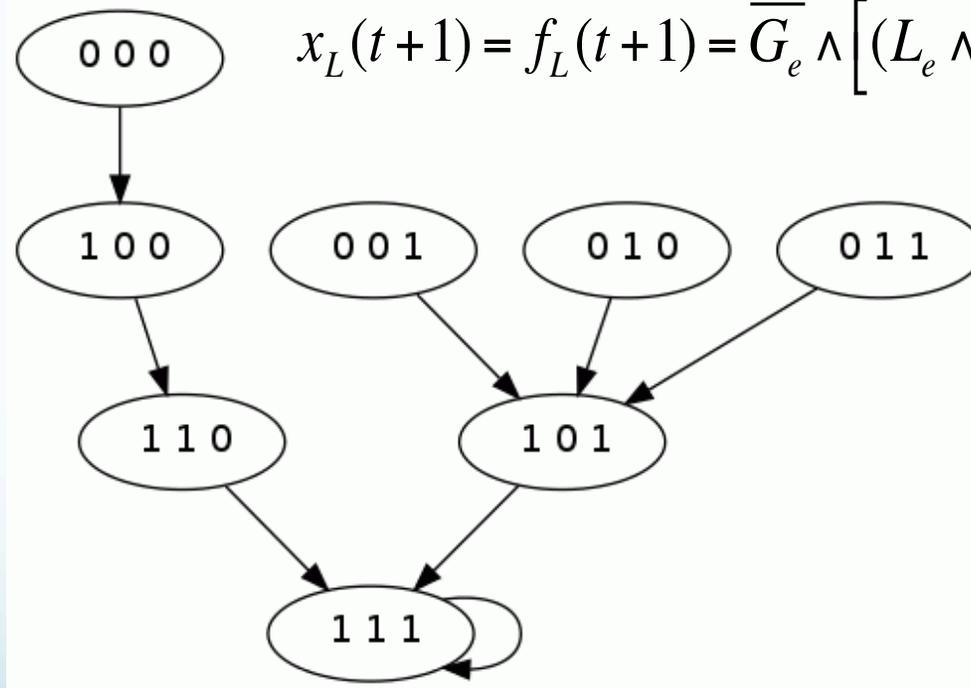
- Simulation of all trajectories (< 20 nodes)

Analysis of Dynamic Algebraic Models (ADAM) v1.1

$$x_M(t+1) = f_M(t+1) = \overline{G_e} \wedge (L(t) \vee L_e)$$

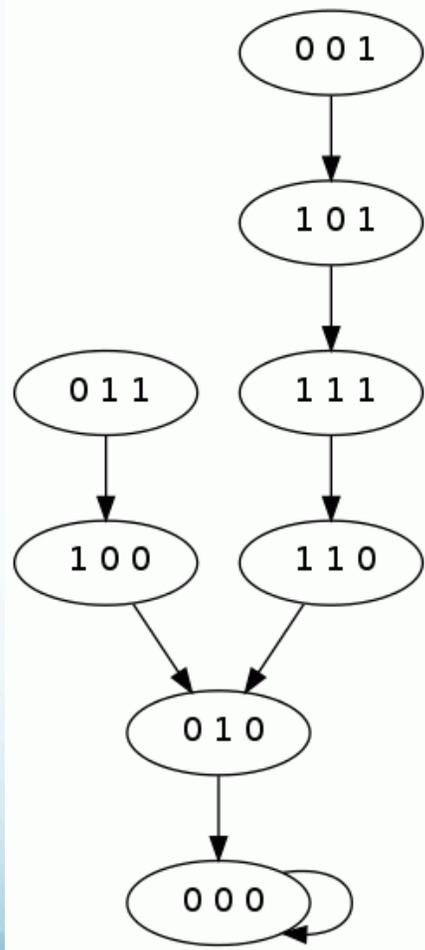
$$x_E(t+1) = f_E(t+1) = M(t)$$

$$x_L(t+1) = f_L(t+1) = \overline{G_e} \wedge [(L_e \wedge E(t)) \vee (L(t) \wedge \overline{E(t)})]$$



State space when $(G_e, L_e) = (0, 1)$. The operon is ON.

Analysis of Dynamic Algebraic Models (ADAM) v1.1



$$x_M(t+1) = f_M(t+1) = \overline{G_e} \wedge (L(t) \vee L_e)$$

$$x_E(t+1) = f_E(t+1) = M(t)$$

$$x_L(t+1) = f_L(t+1) = \overline{G_e} \wedge \left[(L_e \wedge E(t)) \vee (L(t) \wedge \overline{E(t)}) \right]$$

State space when $(G_e, L_e) = (0, 0)$.

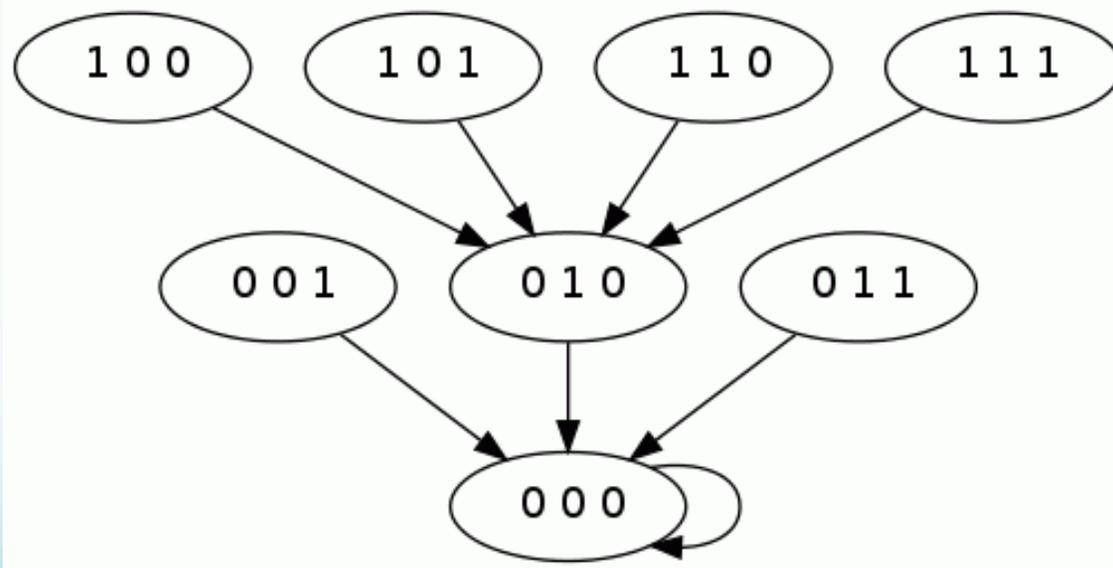
The operon is OFF.

Analysis of Dynamic Algebraic Models (ADAM) v1.1

$$x_M(t+1) = f_M(t+1) = \overline{G_e} \wedge (L(t) \vee L_e)$$

$$x_E(t+1) = f_E(t+1) = M(t)$$

$$x_L(t+1) = f_L(t+1) = \overline{G_e} \wedge [(L_e \wedge E(t)) \vee (L(t) \wedge \overline{E(t)})]$$



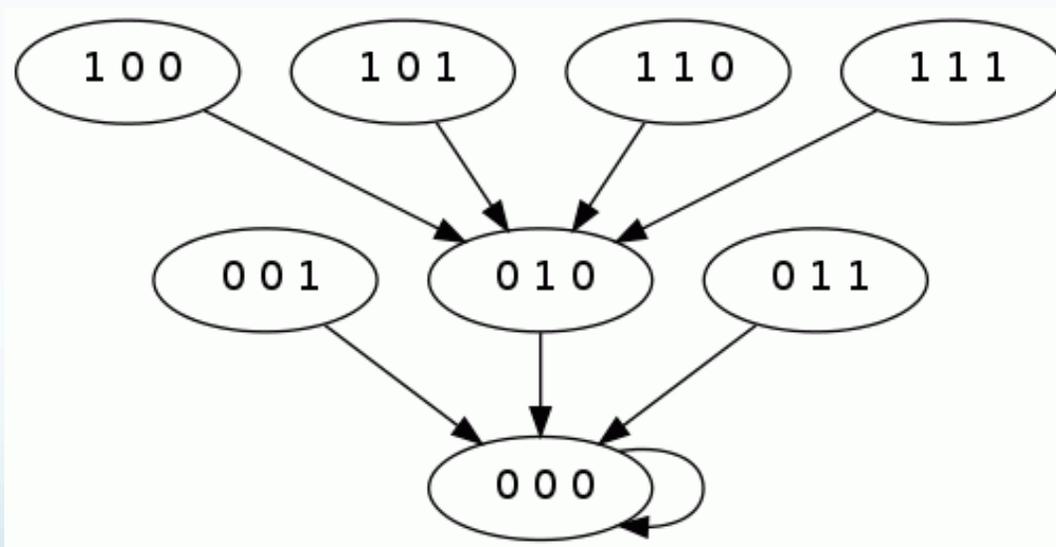
State space when $(G_e, L_e) = (1, 0)$. The operon is OFF.

Analysis of Dynamic Algebraic Models (ADAM) v1.1

$$x_M(t+1) = f_M(t+1) = \overline{G_e} \wedge (L(t) \vee L_e)$$

$$x_E(t+1) = f_E(t+1) = M(t)$$

$$x_L(t+1) = f_L(t+1) = \overline{G_e} \wedge [(L_e \wedge E(t)) \vee (L(t) \wedge \overline{E(t)})]$$



State space when $(G_e, L_e) = (1, 1)$. The operon is OFF.

Summary so far

- **Gene regulatory networks** consist of a collection of gene products that interact with each other to control a specific cell function.
- Classically, these have been modeled quantitatively with differential equations (**continuous models**).
- **Boolean networks** take a different approach. They are **discrete models** that are inherently qualitative.
- The **state space** graph encodes all of the dynamics. The most important features are the **fixed points**, and a necessary step in model validation is to check that they are biologically meaningful.
- The model of the *lac* operon shown here is a “toy model”. Next, we will see more complicated models of the *lac* operon that capture intricate biological features of these systems.
- Modeling with Boolean logic is a relatively new concept, first done in the 1970s. It is a popular research topic in the field of **systems biology**.

A more refined model

- Our model only used 3 variables: mRNA (M), enzymes (E), and lactose (L).

- Let's propose a new model with 5 variables:

- M: mRNA

$$f_M = A$$

- B: β -galactosidase

$$f_B = M$$

- A: allolactose

$$f_A = A \vee (L \wedge B)$$

- L: intracellular lactose

$$f_L = P \vee (L \wedge \bar{B})$$

- P: *lac* permease (transporter protein)

$$f_P = M$$

- Assumptions

- Translation and transcription require one unit of time.
- Protein and mRNA degradation require one unit of time
- Lactose metabolism require one unit of time
- Extracellular lactose is always available.
- Extracellular glucose is always unavailable.

Using ADAM to compute the state space

$$f_M = A$$

$$f_B = M$$

$$f_A = A \vee (L \wedge B)$$

$$f_L = P \vee (L \wedge \bar{B})$$

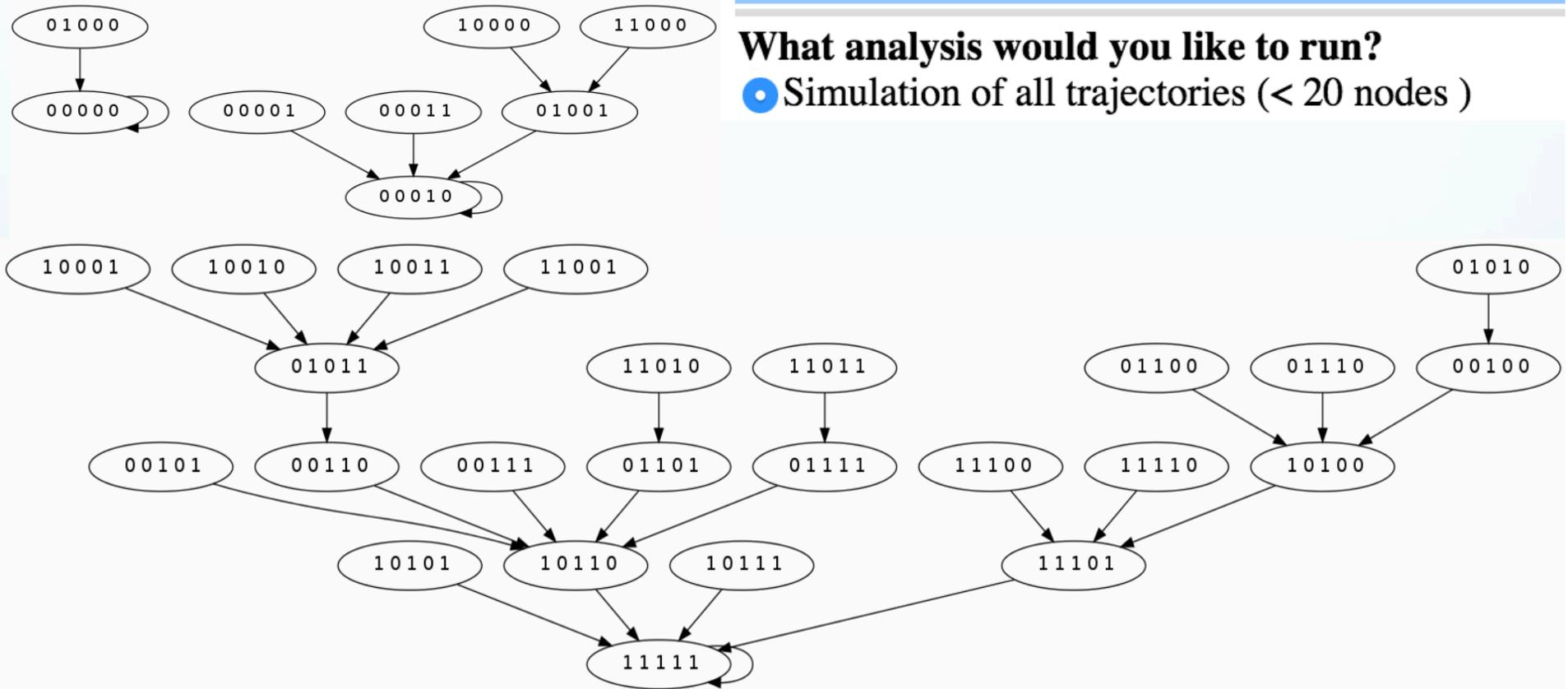
$$f_P = M$$

```
f1 = x3
f2 = x1
f3 = x3+x4*x2+x3*x4*x2
f4 = x5+x4*(1+x2)+x5*x4*(1+x2)
f5 = x1
```

4)

What analysis would you like to run?

- Simulation of all trajectories (< 20 nodes)



Problems with our refined model

- Model variables:

- M: mRNA
- B: β -galactosidase
- A: allolactose
- L: intracellular lactose
- P: *lac* permease (transporter protein)

$$f_M = A$$

$$f_B = M$$

$$f_A = A \vee (L \wedge B)$$

$$f_L = P \vee (L \wedge \bar{B})$$

$$f_P = M$$

- Problems:

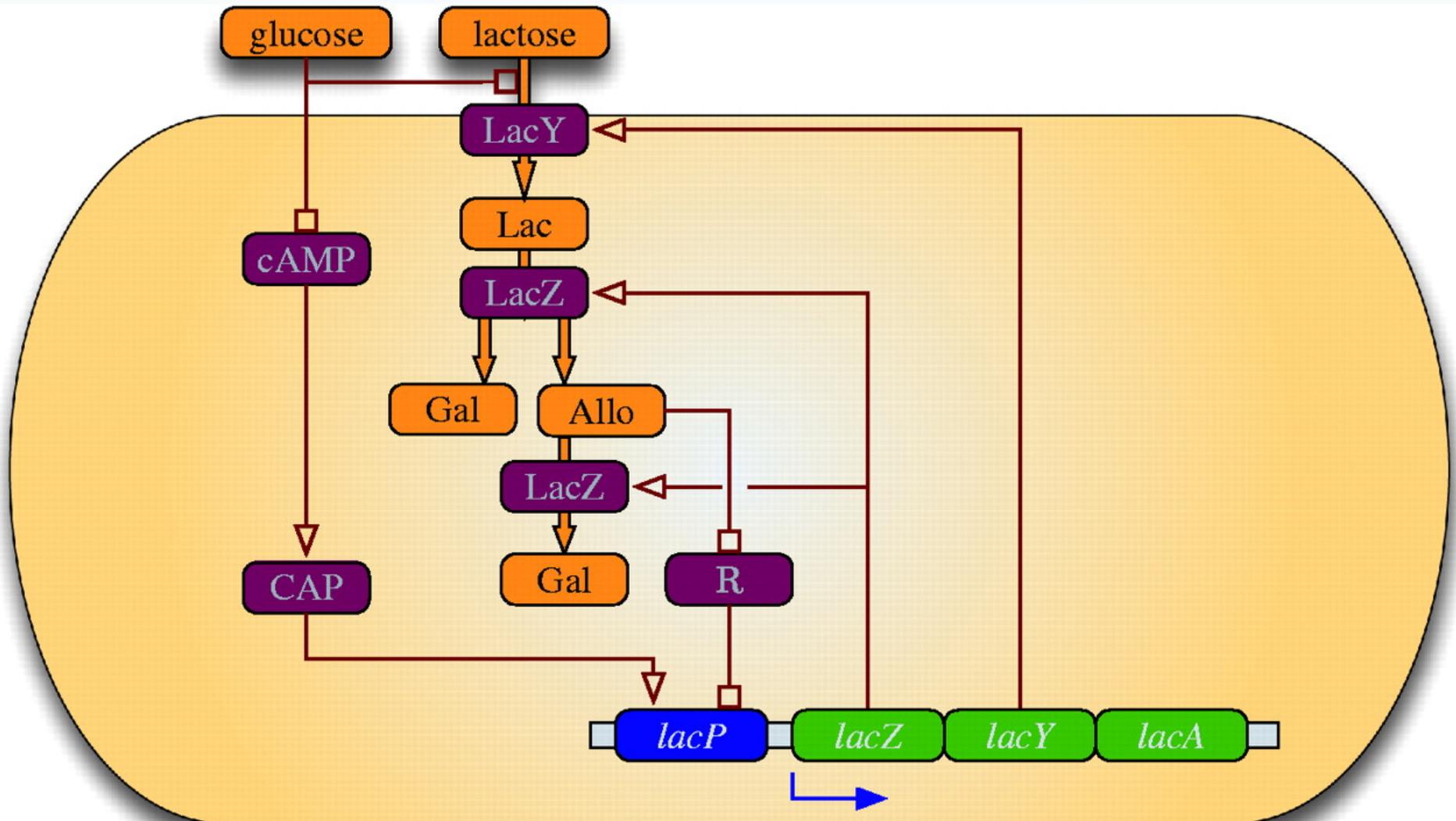
- The fixed point (M,B,A,L,P) = (0,0,0,0,0) should not happen with lactose present but not glucose. [though let's try to justify this...]
- The fixed point (M,B,A,L,P) = (0,0,0,1,0) is not biologically feasible: it would describe a scenario where the bacterium does not metabolize intracellular lactose.

- Conclusion: *The model fails the initial testing and validation, and is in need of modification.* (Homework!)

Catabolite repression

- We haven't yet discussed the cellular mechanism that turns the *lac* operon OFF when both glucose and lactose are present. This is done by **catabolite repression**.
- The *lac* operon promoter region has **2 binding sites**:
 - One for RNA polymerase (this “unzips” and reads the DNA)
 - One for the **CAP-cAMP** complex. This is a complex of two molecules: catabolite activator protein (CAP), and the **cyclic AMP receptor protein** (cAMP, or *crp*).
- Binding of the CAP-cAMP complex is required for transcription for the *lac* operon.
- Intracellular glucose causes the cAMP concentration to decrease.
- When cAMP levels get too low, so do CAP-cAMP complex levels.
- Without the CAP-cAMP complex, the promoter is inactivated, and the *lac* operon is OFF.

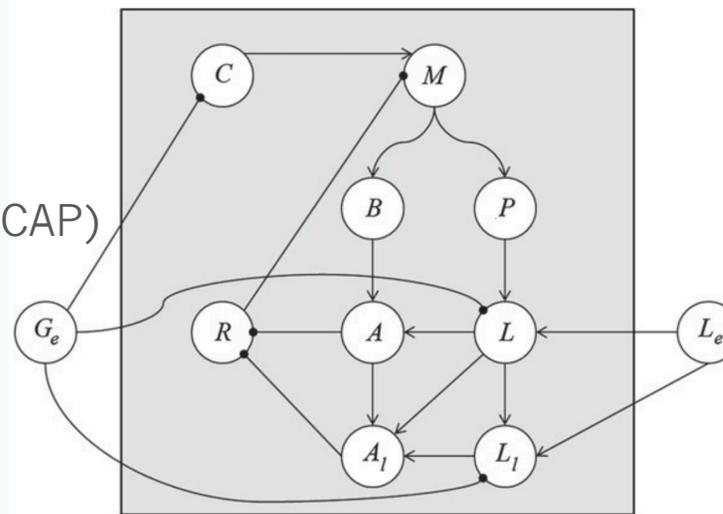
Lac operon gene regulatory network



A more refined model

- Variables:

- M: mRNA
- P: *lac* permease
- B: β -galactosidase
- C: catabolite activator protein (CAP)
- R: repressor protein (LacI)
- A: allolactose
- A_m : at least med. allolactose
- L: intracellular lactose
- L_m : at least med. levels of intracellular lactose



$$f_M = \bar{R} \wedge C$$

$$f_P = M$$

$$f_B = M$$

$$f_C = \bar{G}_e$$

$$f_R = \bar{A} \wedge \bar{A}_m$$

$$f_A = L \wedge B$$

$$f_{A_m} = A \vee L \vee L_m$$

$$f_L = \bar{G}_e \wedge P \wedge L_e$$

$$f_{L_m} = \bar{G}_e \wedge (L \vee L_e)$$

- Assumptions:

- Transcription and translation require 1 unit of time.
- Degradation of all mRNA and proteins occur in 1 time-step.
- High levels of lactose or allolactose at any time t imply (at least) medium levels for the next time-step $t+1$.

A more refined model

- This 9-variable model is about as big as ADAM can render a state space.
- In fact, it doesn't work using the "Open Polynomial Dynamical System (oPDS)" option (variables + parameters).
- Instead, it works under "Polynomial Dynamical System (PDS)", if we manually enter numbers for the parameters.
- Here's a sample piece of the state space:

$$f_M = \bar{R} \wedge C$$

$$f_P = M$$

$$f_B = M$$

$$f_C = \bar{G}_e$$

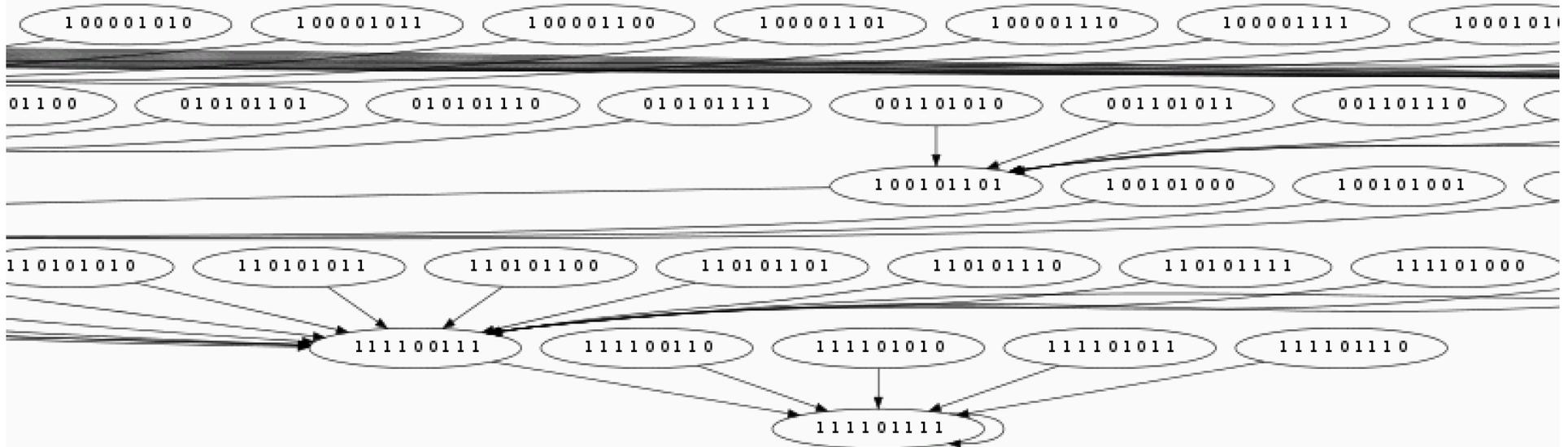
$$f_R = \bar{A} \wedge \bar{A}_m$$

$$f_A = L \wedge B$$

$$f_{A_m} = A \vee L \vee L_m$$

$$f_L = \bar{G}_e \wedge P \wedge L_e$$

$$f_{L_m} = \bar{G}_e \wedge (L \vee L_e)$$



What if the state space is too big?

- The previous 9-variable model is about as big as ADAM can handle.
- However, many gene regulatory networks are much bigger.
 - A Boolean network model (2006) of T helper cell differentiation has 23 nodes, and thus a state space of size $2^{23} = 8,388,608$.
 - A Boolean network model (2003) of the segment polarity genes in *Drosophila melanogaster* (fruit fly) has 60 nodes, and a state space of size $2^{60} \approx 1.15 \times 10^{18}$.
 - There are many more examples...
- For these systems, we need to be able to analyze them without constructing the entire state space.
- Our first goal is to find the fixed points. This amounts to solving a system of equations:

$$\begin{cases} f_{x_1} = x_1 \\ f_{x_2} = x_2 \\ \vdots \\ f_{x_n} = x_n \end{cases}$$

$$f_M = \overline{R} \wedge C$$

$$f_P = M$$

$$f_B = M$$

$$f_C = \overline{G_e}$$

$$f_R = \overline{A} \wedge \overline{A_m}$$

$$f_A = L \wedge B$$

$$f_{A_m} = A \vee L \vee L_m$$

$$f_L = \overline{G_e} \wedge P \wedge L_e$$

$$f_{L_m} = \overline{G_e} \wedge (L \vee L_e)$$

How to find the fixed points

- Let's rename variables: $(M, P, B, C, R, A, A_m, L, L_m) = (x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9)$
- Writing each function in polynomial form, and then $f_{x_i} = x_i$ for each $i=1, \dots, 9$ yields the following system:

$$\begin{array}{l}
 f_M = \bar{R} \wedge C = M \\
 f_P = M = P \\
 f_B = M = B \\
 f_C = \bar{G}_e = C \\
 f_R = \bar{A} \wedge \bar{A}_m = R \\
 f_A = L \wedge B = A \\
 f_{A_m} = A \vee L \vee L_m = A_m \\
 f_L = \bar{G}_e \wedge P \wedge L_e = A_m \\
 f_{L_m} = \bar{G}_e \wedge (L \vee L_e) = L_m
 \end{array}
 \left\{ \begin{array}{l}
 x_1 + x_4 x_5 + x_4 = 0 \\
 x_1 + x_2 = 0 \\
 x_1 + x_3 = 0 \\
 x_4 + (G_e + 1) = 0 \\
 x_5 + x_6 x_7 + x_6 + x_7 + 1 = 0 \\
 x_6 + x_3 x_8 = 0 \\
 x_6 + x_7 + x_8 + x_9 + x_8 x_9 + x_6 x_8 + x_6 x_9 + x_6 x_8 x_9 = 0 \\
 x_8 + x_2 L_e (G_e + 1) = 0 \\
 x_9 + (G_e + 1)(x_8 + x_8 L_e + L_e) = 0
 \end{array} \right.$$

- We need to solve this for all 4 combinations: $(G_e, L_e) = (0, 0), (0, 1), (1, 0), (1, 1)$

How to find the fixed points

- Let's first consider the case when $(G_e, L_e) = (1, 1)$
- We can solve the system by typing the following commands into Sage (<https://cloud.sagemath.com/>), the free open-source mathematical software:

```
1
2 P.<x1,x2,x3,x4,x5,x6,x7,x8,x9> = PolynomialRing(GF(2), 9, order='lex'); P
3     Multivariate Polynomial Ring in x1, x2, x3, x4, x5, x6, x7, x8, x9 over Finite Field of size 2
4
5 Le=1;
6 Ge=1;
7 print "Le =", Le;
8 print "Ge =", Ge;
9
9     Le = 1
10    Ge = 1
11
12 I = ideal(x1+x4*x5+x4, x1+x2, x1+x3, x4+(Ge+1), x5+x6*x7+x6+x7+1, x6+x3*x8,
13 x6+x7+x8+x9+x8*x9+x6*x8+x6*x9+x6*x8*x9, x8+Le*(Ge+1)*x2, x9+(Ge+1)*(Le+x8+Le*x8)); I
14
15 Ideal (x1 + x4*x5 + x4, x1 + x2, x1 + x3, x4, x5 + x6*x7 + x6 + x7 + 1, x3*x8 + x6, x6*x8*x9 +
16 x6*x8 + x6*x9 + x6 + x7 + x8*x9 + x8 + x9, x8, x9) of Multivariate Polynomial Ring in x1, x2,
17 x3, x4, x5, x6, x7, x8, x9 over Finite Field of size 2
18
19 B = I.groebner_basis(); B
20
21 [x1, x2, x3, x4, x5 + 1, x6, x7, x8, x9]
```

What those Sage commands mean

Let's go over what the following commands mean:

- `P.<x1,x2,x3,x4,x5,x6,x7,x8,x9> = PolynomialRing(GF(2),9,order='lex');`
 - Define P to be **the polynomial ring** over 9 variables, x_1, \dots, x_9 .
 - $\text{GF}(2) = \{0,1\}$ because the coefficients are binary.
 - `order='lex'` specifies a **monomial order**. More on this later.

- `Le=1; Ge=1; print "Le =", Le; print "Ge =", Ge;`
 - This defines two constants $(G_e, L_e) = (1, 1)$ and prints them.

- `I = ideal(x1+x4*x5+x4, x1+x2, x1+x3, x4+(Ge+1), x5+x6*x7+x6+x7+1, x6+x3*x8, x6+x7+x8+x9+x8*x9+x6*x8+x6*x9+x6*x8*x9, x8+Le*(Ge+1)*x2, x9+(Ge+1)*(Le+x8+Le*x8)); I`
 - Defines I to be the **ideal** generated by those following 9 polynomials, i.e.,

$$I = \{p_1 f_1 + \dots + p_k f_k : p_k \in P\}$$

- `B = I.groebner_basis(); B`
 - Define B to be the **Gröbner basis** of I w.r.t. the **lex monomial order**. (More on this later)

What does a Gröbner basis tell us?

The output of `B = I.groebner_basis();` B is the following:

$$[x_1, x_2, x_3, x_4, x_5+1, x_6, x_7, x_8, x_9]$$

This is short-hand for the following system of equations:

$$\{x_1 = 0, x_2 = 0, x_3 = 0, x_4 = 0, x_5 + 1 = 0, x_6 = 0, x_7 = 0, x_8 = 0, x_9 = 0\}$$

This simple system has the **same set of solutions** as the much more complicated system we started with:

$$\left\{ \begin{array}{l} x_1 + x_4 x_5 + x_4 = 0 \\ x_1 + x_2 = 0 \\ x_1 + x_3 = 0 \\ x_4 + (G_e + 1) = 0 \\ x_5 + x_6 x_7 + x_6 + x_7 + 1 = 0 \\ x_6 + x_3 x_8 = 0 \\ x_6 + x_7 + x_8 + x_9 + x_8 x_9 + x_6 x_8 + x_6 x_9 + x_6 x_8 x_9 = 0 \\ x_8 + x_2 L_e (G_e + 1) = 0 \\ x_9 + (G_e + 1)(x_8 + x_8 L_e + L_e) = 0 \end{array} \right.$$

Gröbner bases vs. Gaussian elimination

✧ Gröbner bases are a generalization of Gaussian elimination, but for systems of polynomials (instead of systems of linear equations)

✧ In both cases:

- The input is a complicated system that we wish to solve.
- The output is a simple system that we can easily solve by inspection.

✧ Consider the following example:

- Input: The 2x2 system of linear equations
$$\begin{cases} x + 2y = 1 \\ 3x + 8y = 1 \end{cases}$$
- Gaussian elimination yields the following:

$$\left[\begin{array}{cc|c} 1 & 2 & 1 \\ 3 & 8 & 1 \end{array} \right] \rightarrow \left[\begin{array}{cc|c} 1 & 2 & 1 \\ 0 & 2 & -2 \end{array} \right] \rightarrow \left[\begin{array}{cc|c} 1 & 0 & 3 \\ 0 & 2 & -2 \end{array} \right] \rightarrow \left[\begin{array}{cc|c} 1 & 0 & 3 \\ 0 & 1 & -1 \end{array} \right]$$

- This is just the much simpler system with the same solution!
$$\begin{cases} x + 0y = 3 \\ 0x + y = -1 \end{cases}$$

Back-substitution & Gaussian elimination

- ✧ We don't necessarily need to do Gaussian elimination until the matrix is the identity. As long as it is **upper-triangular**, we can back-substitute and solve by hand.

- ✧ For example:

$$\begin{cases} x + z = 2 \\ y - z = 8 \\ 0 = 0 \end{cases}$$

- ✧ Similarly, when Sage outputs a Gröbner basis, it will be in “upper-triangular form”, and we can solve the system easily by back-substituting.
- ✧ We'll do an example right away. For this part of the class, you can think of Gröbner bases as a mysterious “**black box**” that does what we want.
- ✧ We'll study them in more detail shortly, and understand what's going on behind the scenes.

Gröbner bases: an example

✧ Let's use Sage to solve the following system:

$$\begin{cases} x^2 + y^2 + z^2 = 1 \\ x^2 - y + z^2 = 0 \\ x - z = 0 \end{cases}$$

```
17 P.<x,y,z>=PolynomialRing(RR,3,order='lex'); P
18      Multivariate Polynomial Ring in x, y, z over Real Field with 53 bits of precision
19
20 I = ideal(x^2+y^2+z^2-1, x^2-y+z^2, x-z); I
21      Ideal (x^2 + y^2 + z^2 - 1.000000000000000, x^2 - y + z^2, x - z) of Multivariate Polynomial
      Ring in x, y, z over Real Field with 53 bits of precision
22
23 B = I.groebner_basis(); B
24      [x - z, y - 2.000000000000000*z^2, z^4 + 0.500000000000000*z^2 - 0.250000000000000]
```

✧ From this, we get an “upper-triangular” system:

✧ This is something we can solve by hand.

$$\begin{cases} x - z = 0 \\ y - 2z^2 = 0 \\ z^4 + .5z^2 - .25 = 0 \end{cases}$$

Gröbner bases: an example (cont.)

✧ To solve the reduced system:

$$\begin{cases} x - z = 0 \\ y - 2z^2 = 0 \\ z^4 + .5z^2 - .25 = 0 \end{cases}$$

▪ Solve for z in Eq. 3: $z = \pm \sqrt{\frac{-1 + \sqrt{5}}{4}}$

▪ Plug z into Eq. 2 and solve for y : $y = 2z^2 = \frac{-1 + \sqrt{5}}{2}$

▪ Plug y & z into Eq. 1 and solve for x : $x = z = \pm \sqrt{\frac{-1 + \sqrt{5}}{4}}$

$$\begin{cases} x^2 + y^2 + z^2 = 1 \\ x^2 - y + z^2 = 0 \\ x - z = 0 \end{cases}$$

✧ Thus, we get 2 solutions to the original system:

$$(x_1, y_1, z_1) = \left(\sqrt{\frac{-1 + \sqrt{5}}{4}}, \frac{-1 + \sqrt{5}}{2}, \sqrt{\frac{-1 + \sqrt{5}}{4}} \right)$$

$$(x_2, y_2, z_2) = \left(-\sqrt{\frac{-1 + \sqrt{5}}{4}}, \frac{-1 + \sqrt{5}}{2}, -\sqrt{\frac{-1 + \sqrt{5}}{4}} \right)$$

Returning to the *lac* operon

- We have 9 variables: $(M, P, B, C, R, A, A_m, L, L_m) = (x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9)$
- Writing each function in polynomial form, we need to solve the system $f_{x_i} = x_i$ for each $i=1, \dots, 9$, which is the following:

$$\begin{array}{l}
 f_M = \bar{R} \wedge C = M \\
 f_P = M = P \\
 f_B = M = B \\
 f_C = \bar{G}_e = C \\
 f_R = \bar{A} \wedge \bar{A}_m = R \\
 f_A = L \wedge B = A \\
 f_{A_m} = A \vee L \vee L_m = A_m \\
 f_L = \bar{G}_e \wedge P \wedge L_e = A_m \\
 f_{L_m} = \bar{G}_e \wedge (L \vee L_e) = L_m
 \end{array}
 \left\{ \begin{array}{l}
 x_1 + x_4 x_5 + x_4 = 0 \\
 x_1 + x_2 = 0 \\
 x_1 + x_3 = 0 \\
 x_4 + (G_e + 1) = 0 \\
 x_5 + x_6 x_7 + x_6 + x_7 + 1 = 0 \\
 x_6 + x_3 x_8 = 0 \\
 x_6 + x_7 + x_8 + x_9 + x_8 x_9 + x_6 x_8 + x_6 x_9 + x_6 x_8 x_9 = 0 \\
 x_8 + x_2 L_e (G_e + 1) = 0 \\
 x_9 + (G_e + 1)(x_8 + x_8 L_e + L_e) = 0
 \end{array} \right.$$

- We need to solve this for all 4 combinations: $(G_e, L_e) = (0, 0), (0, 1), (1, 0), (1, 1)$ (we already did (1,1)).

Returning to the *lac* operon

- Again, we use variables $(M, P, B, C, R, A, A_m, L, L_m) = (x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9)$ and parameters $(G_e, L_e) = (0, 0)$
- Here is the output from Sage:

```
1
2 P.<x1,x2,x3,x4,x5,x6,x7,x8,x9> = PolynomialRing(GF(2), 9, order = 'lex'); P
3   Multivariate Polynomial Ring in x1, x2, x3, x4, x5, x6, x7, x8, x9 over Finite Field of size 2
4
5 Le=0;
6 Ge=0;
7 print "Le =", Le;
8 print "Ge =", Ge;
9
9   Le = 0
   Ge = 0
10
11 I = ideal(x1+x4*x5+x4, x1+x2, x1+x3, x4+(Ge+1), x5+x6*x7+x6+x7+1, x6+x3*x8,
12 x6+x7+x8+x9+x8*x9+x6*x8+x6*x9+x6*x8*x9, x8+Le*(Ge+1)*x2, x9+(Ge+1)*(Le+x8+Le*x8)); I
12   Ideal (x1 + x4*x5 + x4, x1 + x2, x1 + x3, x4 + 1, x5 + x6*x7 + x6 + x7 + 1, x3*x8 + x6, x6*x8*x9 +
   x6*x8 + x6*x9 + x6 + x7 + x8*x9 + x8 + x9, x8, x8 + x9) of Multivariate Polynomial Ring in x1, x2
   , x3, x4, x5, x6, x7, x8, x9 over Finite Field of size 2
13
14 B = I.groebner_basis(); B
15   [x1, x2, x3, x4 + 1, x5 + 1, x6, x7, x8, x9]
```

- $(M, P, B, C, R, A, A_m, L, L_m) = (x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9) = (0, 0, 0, 1, 1, 0, 0, 0, 0)$

Returning to the *lac* operon

- Again, we use variables $(M, P, B, C, R, A, A_m, L, L_m) = (x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9)$ and parameters $(G_e, L_e) = (1, 0)$
- Here is the output from Sage:

```
1 |
2 | P.<x1,x2,x3,x4,x5,x6,x7,x8,x9> = PolynomialRing(GF(2), 9, order = 'lex'); P
3 |   Multivariate Polynomial Ring in x1, x2, x3, x4, x5, x6, x7, x8, x9 over Finite Field of size 2
4 |
5 | Le=0;
6 | Ge=1;
7 | print "Le =", Le;
8 | print "Ge =", Ge;
9 |
9 |   Le = 0
9 |   Ge = 1
10 |
11 | I = ideal(x1+x4*x5+x4, x1+x2, x1+x3, x4+(Ge+1), x5+x6*x7+x6+x7+1, x6+x3*x8,
11 | x6+x7+x8+x9+x8*x9+x6*x8+x6*x9+x6*x8*x9, x8+Le*(Ge+1)*x2, x9+(Ge+1)*(Le+x8+Le*x8)); I
12 |
12 |   Ideal (x1 + x4*x5 + x4, x1 + x2, x1 + x3, x4, x5 + x6*x7 + x6 + x7 + 1, x3*x8 + x6, x6*x8*x9 +
12 | x6*x8 + x6*x9 + x6 + x7 + x8*x9 + x8 + x9, x8, x9) of Multivariate Polynomial Ring in x1, x2,
12 | x3, x4, x5, x6, x7, x8, x9 over Finite Field of size 2
13 |
14 | B = I.groebner_basis(); B
15 |
15 | [x1, x2, x3, x4, x5 + 1, x6, x7, x8, x9]
```

- $(M, P, B, C, R, A, A_m, L, L_m) = (x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9) = (0, 0, 0, 0, 1, 0, 0, 0, 0)$

Returning to the *lac* operon

- Again, we use variables $(M, P, B, C, R, A, A_m, L, L_m) = (x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9)$ and parameters $(G_e, L_e) = (0, 1)$
- Here is the output from Sage:

```
1
2 P.<x1,x2,x3,x4,x5,x6,x7,x8,x9> = PolynomialRing(GF(2), 9, order = 'lex'); P
3   Multivariate Polynomial Ring in x1, x2, x3, x4, x5, x6, x7, x8, x9 over Finite Field of size 2
4
5 Le=0;
6 Ge=1;
7 print "Le =", Le;
8 print "Ge =", Ge;
9
10   Le = 0
11   Ge = 1
12
13 I = ideal(x1+x4*x5+x4, x1+x2, x1+x3, x4+(Ge+1), x5+x6*x7+x6+x7+1, x6+x3*x8,
14 x6+x7+x8+x9+x8*x9+x6*x8+x6*x9+x6*x8*x9, x8+Le*(Ge+1)*x2, x9+(Ge+1)*(Le+x8+Le*x8)); I
15
16   Ideal (x1 + x4*x5 + x4, x1 + x2, x1 + x3, x4, x5 + x6*x7 + x6 + x7 + 1, x3*x8 + x6, x6*x8*x9 +
17   x6*x8 + x6*x9 + x6 + x7 + x8*x9 + x8 + x9, x8, x9) of Multivariate Polynomial Ring in x1, x2,
18   x3, x4, x5, x6, x7, x8, x9 over Finite Field of size 2
19
20 B = I.groebner_basis(); B
21
22   [x1, x2, x3, x4, x5 + 1, x6, x7, x8, x9]
```

$$(M, P, B, C, R, A, A_m, L, L_m) = (x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9) = (1, 1, 1, 1, 0, 1, 1, 1, 1)$$

Fixed point analysis of the *lac* operon

Using the variables $(M, P, B, C, R, A, A_m, L, L_m) = (x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9)$

we got the following fixed points for each choice of parameters (G_e, L_e)

- Input: $(G_e, L_e) = (0, 0)$
Fixed point: $(x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9) = (0, 0, 0, 1, 1, 0, 0, 0, 0)$
- Input: $(G_e, L_e) = (1, 0)$
Fixed point: $(x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9) = (0, 0, 0, 0, 1, 0, 0, 0, 0)$
- Input: $(G_e, L_e) = (1, 1)$
Fixed point: $(x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9) = (0, 0, 0, 0, 1, 0, 0, 0, 0)$
- Input: $(G_e, L_e) = (0, 1)$
Fixed point: $(x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9) = (1, 1, 1, 1, 0, 1, 1, 1, 1)$

All of these fixed points make biological sense!