# Lecture 5.1: Basic cryptographic ciphers 

Matthew Macauley<br>Department of Mathematical Sciences<br>Clemson University<br>http://www.math.clemson.edu/~macaule/

Math 4190, Discrete Mathematical Structures

## Encoding messages as numbers

In this lecture, we'll see how to send encoded messages, which will be numbers.
We can encode any word as a number in base-26:

| A | B | C | D | E | F | G | H | I | J | K | L | M | N | O | P | Q | R | S | T | U | V | W | X | Y | Z |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 | 21 | 22 | 23 | 24 | 25 |

In base 26, the word CLEMSON can be encoded as $2 \begin{array}{llllll}11 & 4 & 12 & 18 & 14 & 13 .\end{array}$
We can convert this to decimal (base 10):

$$
2 \cdot 26^{0}+11 \cdot 26^{1}+4 \cdot 26^{2}+12 \cdot 26^{3}+18 \cdot 26^{4}+14 \cdot 26^{5}+13 \cdot 26^{6}=4190683824
$$

To reverse this process, recursively divide 26 into the number. Let's try this with 221707947:

| 221707947 | $=$ | $8527228 \cdot 26$ | +19 | $\mathbf{T}$ |
| ---: | ---: | ---: | ---: | ---: |
| 8527228 | $=$ | $327970 \cdot 26$ | +8 | $\mathbf{I}$ |
| 327970 | $=$ | $12614 \cdot 26$ | + | 6 |
| $\mathbf{G}$ |  |  |  |  |
| 12614 | $=$ | $485 \cdot 26$ | + | 4 |
| $\mathbf{E}$ |  |  |  |  |
| 485 |  | $18 \cdot 26$ | +17 | $\mathbf{R}$ |
| 18 |  | $0 \cdot 26$ | +18 | $\mathbf{S}$ |

Now, suppose that we wanted to send this as a secret message...

## Some history

Though he wasn't the first, Julius Caesar (100 B.C-44 B.C) used an encryption device called a cipher in his private correspondences.

An encrypted message would looks something like this: RZ WKDWV VKDUS


Decrypted message: OW THATS SHARP

## Caesar cipher

The Caesar cipher is defined by the following:

- key, $k \in N$,
- encryption function, $e(x)$,
- decryption function, $d(x)$,

$$
e(x)=x+k \quad(\bmod 26), \quad d(y)=y-k \quad(\bmod 26)
$$

We first associate each letter to a number in $\mathbb{Z}_{26}:=\{0,1, \ldots, 24,25\}$, as follows:

| A | B | C | D | E | F | G | H | I | J | K | L | M | N | O | P | Q | R | S | T | U | V | W | X | Y | Z |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 | 21 | 22 | 23 | 24 | 25 |

To see an example of this, suppose that $k=18$.

We encrypt the letter $\mathbf{R}$ by

$$
\begin{aligned}
e(17) & \equiv 17+18 \quad(\bmod 26) \\
& \equiv 35 \quad(\bmod 26) \\
& \equiv 9 \quad(\bmod 26)
\end{aligned}
$$

which is $\mathbf{J}$.

Let's decrypt L:

$$
\begin{aligned}
d(11) & \equiv 11-18 \quad(\bmod 26) \\
& \equiv-7 \quad(\bmod 26) \\
& \equiv 19 \quad(\bmod 26)
\end{aligned}
$$

which is $\mathbf{T}$.

## A cipher with multiplication

Consider the following encryption function:

$$
e: \mathbb{Z}_{26} \rightarrow \mathbb{Z}_{26}, \quad e(x)=5 x \quad(\bmod 26)
$$

This works because the function $e$ is injective, and this is because $\operatorname{gcd}(26,5)=1$.
Let's encrypt the letter $R$, which is $x=17$ :

$$
\begin{aligned}
e(17) & \equiv 5 \cdot 17 \quad(\bmod 26) \\
& \equiv 85 \quad(\bmod 26) \\
& \equiv 7 \quad(\bmod 26)
\end{aligned}
$$

The decryption function is

$$
d(x)=21 x \quad(\bmod 26)
$$

This works because in $\mathbb{Z}_{26}$, the multiplicative inverse of $k=5$ is $5^{-1}:=21$ :

$$
\begin{aligned}
5 \cdot 21 & \equiv 105 \quad(\bmod 26) \\
& \equiv 1 \quad(\bmod 26)
\end{aligned}
$$

## Number theory fact

A number $k \in \mathbb{Z}_{n}$ has a multiplicative inverse iff $\operatorname{gcd}(n, k)=1$.

## A cipher with multiplication and addition

Consider the following encrypting function:

$$
e: \mathbb{Z}_{26} \rightarrow \mathbb{Z}_{26}, \quad e(x)=5 x+3 \quad(\bmod 26)
$$

In other words, given the input $x$, we:

1. multiply by 5
2. add 3.

To decrypt a message $m$, we need to "undo these" in the opposite order:
2. subtract 3

1. multiply by $5^{-1}=21$.

The decryption function is thus

$$
d(x)=21(x-3) \quad(\bmod 26)
$$

## A weakness of character ciphers

The ciphers that we've seen are called character, or monographic ciphers: all copies of the same letter get encrypted the same way:

$$
e\left(x_{i}\right)=e\left(x_{j}\right) \Rightarrow x_{i}=x_{j}
$$

If the message is long, the the private key can be deduced by analyzing letter frequencies.


## Block ciphers

A more sophisticated cipher are the block, or polygraphic ciphers, which encrypt blocks of plaintext letters to blocks of ciphertext letters of the same length.

One such system was developed by Blaise de Vigenère in 1585, called the Vigenère cipher.
We'll introduce this by an example.

$$
\begin{array}{llllllllllllllllllllllllll}
\text { A } & \text { B } & \text { C } & \text { D } & \text { E } & \text { F } & \text { G } & \text { H } & \text { I } & \text { J } & \text { K } & \text { L } & \text { M } & \text { N } & \text { O } & \text { P } & \text { Q } & \text { R } & \text { S } & \text { T } & \text { U } & \text { V } & \text { W } & \text { X } & \text { Y } & \text { Z } \\
0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 & 12 & 13 & 14 & 15 & 16 & 17 & 18 & 19 & 20 & 21 & 22 & 23 & 24 & 25
\end{array}
$$

Let's encrypt the message ENGINEERING:

$$
p_{1} p_{2} p_{3} p_{4} p_{5} p_{6} p_{7} p_{8} p_{9} p_{10} p_{11}=\begin{array}{lllllllllll}
4 & 13 & 6 & 8 & 13 & 4 & 4 & 17 & 8 & 13 & 6
\end{array}
$$

using the key ROCKS:

$$
k_{1} k_{2} k_{3} k_{4} k_{5}=17 \quad 14 \quad 2 \quad 1018 .
$$

$$
c_{i}=\begin{array}{c|ccccc|ccccc|c} 
\\
p_{i} \\
k_{i} & \mathbf{E} & \mathbf{N} & \mathbf{G} & \mathbf{I} & \mathbf{N} & \mathbf{E} & \mathbf{E} & \mathbf{R} & \mathbf{I} & \mathbf{N} & \mathbf{G} \\
4 & 13 & 6 & 8 & 13 & 4 & 4 & 17 & 8 & 13 & 6 \\
17 & 14 & 2 & 10 & 18 & 17 & 14 & 2 & 10 & 18 & 17 \\
21 & 1 & 8 & 18 & 5 & 21 & 18 & 19 & 18 & 5 & 23 \\
\mathbf{V} & \mathbf{B} & \mathbf{I} & \mathbf{S} & \mathbf{F} & \mathbf{V} & \mathbf{S} & \mathbf{T} & \mathbf{S} & \mathbf{F} & \mathbf{X}
\end{array}
$$

## Decryption with the Vigenère cipher

```
A B C D E F G H I J K L M N O P Q R S T U V W X Y Z
0
```

Let's decrypt the message TZGWK FBVSY WFU:

$$
c_{1} c_{2} c_{3} c_{4} c_{5} c_{6} c_{7} c_{8} c_{9} c_{10} c_{11} c_{12} c_{13}=\begin{array}{lllllllllllll}
19 & 25 & 6 & 22 & 10 & 5 & 1 & 21 & 18 & 24 & 22 & 5 & 20
\end{array}
$$

using the same key ROCKS:

$$
k_{1} k_{2} k_{3} k_{4} k_{5}=17 \quad 14 \quad 2 \quad 1018
$$

$$
p_{i}=\begin{array}{c|ccccc|ccccc|ccc} 
\\
c_{i} \\
k_{i} & \mathbf{T} & \mathbf{Z} & \mathbf{G} & \mathbf{W} & \mathbf{K} & \mathbf{F} & \mathbf{B} & \mathbf{V} & \mathbf{S} & \mathbf{Y} & \mathbf{W} & \mathbf{F} & \mathbf{U} \\
c_{i}-k_{i} & 19 & 25 & 6 & 22 & 10 & 5 & 1 & 21 & 18 & 24 & 22 & 5 & 20 \\
17 & 14 & 2 & 10 & 18 & 17 & 14 & 2 & 10 & 18 & 17 & 14 & 2 \\
2 & 11 & 4 & 12 & 18 & 14 & 13 & 19 & 8 & 6 & 4 & 17 & 18 \\
\mathbf{C} & \mathbf{L} & \mathbf{E} & \mathbf{M} & \mathbf{S} & \mathbf{O} & \mathbf{N} & \mathbf{T} & \mathbf{I} & \mathbf{G} & \mathbf{E} & \mathbf{R} & \mathbf{S}
\end{array}
$$

## Different types of ciphers

The ciphers in this lecture are symmetric: Decryption is the opposite ("inverse") of encryption.

Not only is the same key is used for encryption and decryption, but that key needs to be kept private.

## Definition

In an asymmetric cipher, there are two distinct keys:

- A public key, used for encryption;
- A private key, used for decryption.

The instructions for encrypting a meesage can be made public without compromising the security.

